

# Outage Constrained Robust Transmission Design for IRS-aided Secure Communications with Direct Communication Links

Sheng Hong, Cunhua Pan, Gui Zhou, Hong Ren, and Kezhi Wang

**Abstract**—This paper considers an intelligent reflecting surface (IRS) aided secure communication with direct communication links where a legitimate receiver (Bob) served by a base station (BS) is overheard by multiple eavesdroppers (Eves), meanwhile the artificial noise (AN) is incorporated to confuse Eves. Since Eves are not legitimate users, their channels cannot be estimated perfectly. We investigate two scenarios with partial channel state information (CSI) error of only cascaded BS-IRS-Eve channel and full CSI errors of both cascaded BS-IRS-Eve channel and direct BS-Eve channel under the statistical CSI error model. To ensure the security performance under CSI errors, the transmit beamforming, AN spatial distribution at the BS, and phase shifts at IRS are jointly optimized to minimize the transmit power constrained by the minimum data rate requirement of Bob and the outage probability of maximum data rate limitation of Eves. In contrast to existing works, the direct link considered in our work makes the optimization of phase shifts at IRS much more challenging, thus we propose a series of novel and artful mathematical manipulations to tackle this issue. Moreover, the proposed algorithm can be applied for both uncorrelated and correlated CSI errors. Simulations confirm the superiority of our proposed algorithm.

**Index Terms**—Intelligent Reflecting Surface (IRS), CSI errors, robust transmission design, secure communication, outage probability, direct communication link.

## I. INTRODUCTION

THE intelligent reflecting surface (IRS) is a promising technique in future six generation (6G) communication networks [1], [2]. The IRS consists of a large number of reflecting units, each of which can reflect the incident signal

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passively [3]. By properly tuning the phase shifts of reflecting units, the reflected signals can be added constructively or destructively [1]. Thus an IRS can intelligently configure the wireless environment to help the transmissions between the sender and the receiver. Since the reflecting unit of the IRS operates in a passive mode without any radio frequency (RF) chain, deploying an IRS costs much less than deploying a relay [4]. Moreover, it can be easily integrated into traditional communication systems with only minor modifications. Therefore, the IRS-aided wireless communications have received extensive attention in multicell networks [5], mobile edge computing [6], multigroup multicast communication [7], cognitive radio system [8], wireless power transfer design [9], etc.

In view of the great potential, the IRS has recently been exploited to enhance the physical layer security in wireless communications [10]–[12]. Especially in some tough scenarios, where the non-zero secrecy rate is difficult to achieve without the IRS, the profits brought by IRS are more prominent. For example, in [13], an IRS was applied to tackle the challenging scenarios where the channel of the legitimate communication link and that of the eavesdropping link were highly correlated. A scenario where the eavesdroppers (Eves) were closer to the base station (BS) than the desired users was investigated by employing an IRS in [14], [15]. In IRS-aided secure systems, the authors in [16] confirmed the advantages of using artificial noise (AN) to enhance the secrecy rate. The work in [17] further verified the performance enhancement by jointly using an IRS and AN when the direct BS-user links were blocked.

However, all above-mentioned works assumed that the channel state information (CSI) associated with all involved channels is perfectly known at the BS, which is too idealistic. Since Eves are usually unregistered users and the IRS is passive, the Eves' channel, especially the IRS-related channel cannot be perfect. Generally, in IRS-aided systems, the direct channels from BS to users are first estimated by turning off the IRS [18], then the IRS is turned on, and the IRS-related channels are estimated. The direct channels can be estimated by traditional algorithms [18]. As for the IRS-related channels, currently, there are roughly two kinds of estimating approaches: 1) separately estimating the channels of BS-IRS link and IRS-user link [19]. The main idea is to install some active channel estimators at the IRS, then estimate the two separate channels individually, and finally send the estimated channels to the BS; 2) directly estimating the cascaded channels [20]–[23], which is the composite channel of the BS-IRS link and IRS-user link, and can be exploited to achieve the optimal beamforming

design [7], [24]. The second channel estimation (CE) approach is more attractive than the first one, because no additional active hardware is required, no additional power is consumed, and the channel training and feedback overhead is reduced.

In spite of these CE approaches, the CSI errors for Eves are still unavoidable since Eves are not legitimate users. To reduce the performance loss brought by CSI errors, a few recent works have addressed the robust design problem in IRS-aided communications. The earliest research on robust design for IRS-aided communications relies on the first CE approach. A worst-case robust design was investigated in [25] with CSI errors on IRS-user channels in a multiuser MISO communication system. Then, a robust design algorithm in IRS-aided secure communications was proposed in [26], where the IRS-Eve link had bounded CSI errors, and the direct communication links were blocked. Recently, the research on robust design mainly focuses on the the second CE approach, which is more appealing in practice. The robust design based on the second approach was firstly proposed in [27] for IRS-aided MISO communications, where the cascaded BS-IRS-user channels were imperfect with both the bounded and statistical channel errors. Similarly, the robust design for an IRS-aided cognitive radio system was investigated with imperfect cascaded CSI on primary user (PU)-related channels in [28]. For the bounded CSI errors, a worst-case robust transmit power minimization problem was investigated for the IRS-aided green MISO communications based on the second CE method [29], and a robust sum-rate maximization problem was solved for multiuser MISO systems with self-sustainable IRS [30]. For the statistical CSI errors, a robust probabilistic-constrained transmit power minimization problem was investigated for IRS-aided MISO communications without direct links [31]. The outage-constrained robust beamforming problem was transformed into an outage probability minimization problem in [32]. We investigated the robust transmission design in IRS-assisted secure communications in [33] by assuming that the cascaded BS-IRS-Eve channels have statistic CSI errors, and the direct links are blocked.

From above research, we find that the direct communication links were usually assumed to be blocked, e.g., the methods in [31], [33]. Actually, the existence of direct links will make the robust design problem much more challenging, especially for the optimization of IRS phase shifts. Thus, in this paper, we investigate an outage constrained robust power minimization (OCR-PM) problem in an IRS-aided secure communication system with direct links. By considering the partial and full CSI errors on Eves' channels, the transmit beamformer, the AN spatial distribution, and the phase shifts at the IRS are jointly optimized. Moreover, in contrast to existing robust designs under statistical CSI errors which cannot apply to the case with correlated CSI errors, e.g., the methods in [27], [28], we propose an outage constrained robust (OCR) design framework catering for both uncorrelated or correlated CSI errors in IRS-aided secure communications with direct link. The contributions of this work are summarized as follows:

- 1) In contrast to existing literature [26], [33] which assumes that the direct BS-users links were blocked, we investigate the robust design in a more general

and practical scenario with both direct links and IRS reflecting links. The existence of direct communication link makes the optimization of phase shifts at the IRS much more challenging since almost all constraints become nonconvex. To address this issue, we propose a series of artful mathematical manipulations to transform these constraints into convex ones, thus provide a robust design framework for the case with direct link.

- 2) In the robust design problem, we consider two kinds of scenarios where only the CSI errors exist for cascaded channels (i.e. partial CSI errors) as well as for both the cascaded channels and direct channels (i.e. full CSI errors). To solve the formulated problem, the alternation optimization (AO) strategy is leveraged to decouple the optimization variables. The Bernstein-type inequality (BTI) [34] is utilized to tackle the probability constraints. The penalty convex-concave procedure (CCP) [27] is explored to handle the nonconvex unit modulus constraints of IRS phase shifts.
- 3) In contrast to existing research [27], [28] where the covariance matrix of CSI errors is simplified into an identity matrix to facilitate the algorithm design, our proposed algorithm can be applied with more general forms of covariance matrices, which can describe both uncorrelated and correlated CSI errors.
- 4) Simulation results verify the effectiveness of our proposed algorithm under both scenarios of partial CSI errors and full CSI errors, and reveal that the proposed algorithm outperforms the maximum ratio transmission (MRT) and isotropic AN based baseline schemes.

*Notations:* Boldface lowercase and uppercase letters are used to represent vectors and matrices, respectively. The superscripts  $(\cdot)^T$ ,  $(\cdot)^H$ , and  $(\cdot)^*$  stand for the transpose, Hermitian, and conjugate operators, respectively.  $\mathbb{C}^{M \times N}$  represents the set of all  $M \times N$  complex matrices, and  $\mathbb{H}^n$  denotes the set of all  $n \times n$  Hermitian matrices.  $\text{Tr}(\cdot)$ ,  $\text{Re}\{\cdot\}$ , and  $\text{diag}\{\cdot\}$  denote the trace, the real part of a complex value, and a diagonal matrix.  $\text{vec}(\mathbf{A})$  denotes the vectorization operation on the matrix  $\mathbf{A}$ .  $\text{unvec}[\mathbf{a}]_{M \times N}$  means reshaping the vector  $\mathbf{a}$  into an  $M \times N$  dimensional matrix.  $\sigma_{\max}\{\mathbf{A}\}$  represents the maximum singular value of matrix  $\mathbf{A}$ .  $\lambda_{\max}\{\mathbf{A}\}$  and  $\lambda_{\min}\{\mathbf{A}\}$  denote the maximum and minimum eigen value of matrix  $\mathbf{A}$ , respectively.  $\mathcal{CN}(\boldsymbol{\mu}, \mathbf{Z})$  represents a circularly symmetric complex gaussian (CSCG) distribution with a mean vector  $\boldsymbol{\mu}$  and covariance matrix  $\mathbf{Z}$ .

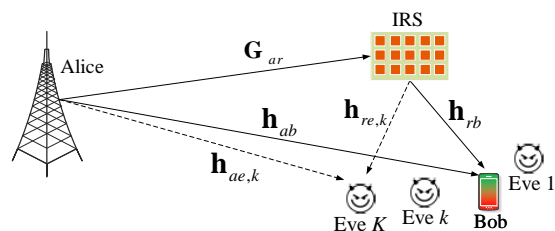


Fig. 1. An IRS-aided secure communication with one Bob and multiple Eves under MISO wiretap channels.

## II. SYSTEM MODEL

### A. Signal Transmission Model

As illustrated in Fig. 1, we consider a wireless downlink scenario, where a single-antenna legitimate receiver (referred as Bob), is overheard by multiple single-antenna eavesdroppers (referred as Eves). Eves are assumed to be participating users (but not legitimate users), thus the transmitter (referred as Alice) has Eves' channel state information (CSI) to some extent, which is not perfect. The Alice is equipped with  $N_t$  antennas, and the IRS is equipped with  $M$  reflection units, thus the spatial degrees of freedom (DoFs) at the transmitter and at the IRS are utilized to degrade the Eves' interceptions. By considering the AN-aided transmit beamforming, the transmit signal vector at Alice can be represented as

$$\mathbf{x} = \mathbf{s} + \mathbf{z} = \mathbf{w}s + \mathbf{z}, \quad (1)$$

where  $s$  is the data symbol intended for Bob, and  $\mathbf{z}$  is the AN generated by Alice to confuse Eves.  $\mathbf{w} \in \mathbb{C}^{N_t \times 1}$  is the transmit beamforming vector. We assume that the transmit signal vector  $\mathbf{s}$  and the noise vector  $\mathbf{z}$  follow the complex Gaussian distributions of  $CN(\mathbf{0}, \mathbf{W})$  and  $CN(\mathbf{0}, \mathbf{Z})$ , respectively, where  $\mathbf{W} = \mathbf{w}\mathbf{w}^H$ . Obviously, both  $\mathbf{W}$  and  $\mathbf{Z}$  are positive semidefinite matrices, and we have  $\text{rank}(\mathbf{W}) = 1$ .

Then the received signals at Bob and the  $k$ th Eve can be respectively expressed as

$$y_b = \hat{\mathbf{h}}_b^H \mathbf{x} + n_b = (\mathbf{h}_{ab}^H + \mathbf{h}_{rb}^H \Phi \mathbf{G}_{ar}) \mathbf{x} + n_b, \quad (2a)$$

$$y_{e,k} = \hat{\mathbf{h}}_{e,k}^H \mathbf{x} + n_{e,k} = (\mathbf{h}_{ae,k}^H + \mathbf{h}_{re,k}^H \Phi \mathbf{G}_{ar}) \mathbf{x} + n_{e,k}, \quad (2b)$$

where  $\forall k \in \mathcal{K} = \{1, 2, \dots, K\}$ . The channels of direct communication links are  $\mathbf{h}_{ab} \in \mathbb{C}^{N_t \times 1}$  and  $\mathbf{h}_{ae,k} \in \mathbb{C}^{N_t \times 1}$ , which respectively denote the Alice-Bob link and Alice-Eve link. The IRS can provide reflecting links to enhance the communication for Bob, and interfere the communication for Eves. The channel from Alice to IRS is modeled by  $\mathbf{G}_{ar} \in \mathbb{C}^{M \times N_t}$ . The channels from IRS to Bob and from IRS to Eves are modeled by  $\mathbf{h}_{rb} \in \mathbb{C}^{M \times 1}$  and  $\mathbf{h}_{re,k} \in \mathbb{C}^{M \times 1}$ , respectively. Let  $\phi_m = e^{j\theta_m}$  and define the reflection coefficients matrix of the IRS by  $\Phi = \text{diag}\{\phi_1, \dots, \phi_m, \dots, \phi_M\}$ , where  $\theta_m \in [0, 2\pi)$  denotes the phase shift of the  $m$ th unit of the IRS. Then the equivalent channel of the composite Alice-IRS-Bob link can be defined by  $\hat{\mathbf{h}}_b \triangleq \mathbf{h}_{ab} + \mathbf{G}_{ar}^H \Phi \mathbf{h}_{rb}$ ,  $\hat{\mathbf{h}}_b \in \mathbb{C}^{N_t \times 1}$ , while the equivalent channel of the composite Alice-IRS-Eve link can be defined by  $\hat{\mathbf{h}}_{e,k} \triangleq \mathbf{h}_{ae,k} + \mathbf{G}_{ar}^H \Phi \mathbf{h}_{re,k}$ ,  $\hat{\mathbf{h}}_{e,k} \in \mathbb{C}^{N_t \times 1}$ .  $n_b \sim CN(0, \sigma_b^2)$  denotes the additive white Gaussian noise (AWGN) received at Bob, while  $n_{e,k} \sim CN(0, \sigma_{e,k}^2)$  denotes the AWGN received at the  $k$ th Eve.

By defining the vector  $\phi = [\phi_1, \dots, \phi_m, \dots, \phi_M]^T$ , the equivalent channel  $\hat{\mathbf{h}}_b$  and  $\hat{\mathbf{h}}_{e,k}$  can be reexpressed as

$$\hat{\mathbf{h}}_b^H \triangleq \mathbf{h}_{ab}^H + \phi^T \mathbf{G}_{cb}, \quad (3a)$$

$$\hat{\mathbf{h}}_{e,k}^H \triangleq \mathbf{h}_{ae,k}^H + \phi^T \mathbf{G}_{ce,k}, \quad \forall k \in \mathcal{K}, \quad (3b)$$

where the  $\mathbf{G}_{cb} \triangleq \text{diag}(\mathbf{h}_{rb}^H) \mathbf{G}_{ar} \in \mathbb{C}^{M \times N_t}$  is defined as the cascaded Alice-IRS-Bob channel, and  $\mathbf{G}_{ce,k} \triangleq \text{diag}(\mathbf{h}_{re,k}^H) \mathbf{G}_{ar} \in \mathbb{C}^{M \times N_t}$  is defined as the cascaded Alice-IRS-Eve channel.

Based on above channel models, the achievable data rates in (bit/s/Hz) of Bob and the  $k$ th Eve are

$$C_b(\mathbf{W}, \mathbf{Z}, \Phi) = \log \left( 1 + \frac{\hat{\mathbf{h}}_b^H \mathbf{W} \hat{\mathbf{h}}_b}{\sigma_b^2 + \hat{\mathbf{h}}_b^H \mathbf{Z} \hat{\mathbf{h}}_b} \right), \quad (4a)$$

$$C_{e,k}(\mathbf{W}, \mathbf{Z}, \Phi) = \log \left( 1 + \frac{\hat{\mathbf{h}}_{e,k}^H \mathbf{W} \hat{\mathbf{h}}_{e,k}}{\sigma_{e,k}^2 + \hat{\mathbf{h}}_{e,k}^H \mathbf{Z} \hat{\mathbf{h}}_{e,k}} \right). \quad (4b)$$

Then the achievable secrecy rate [35] can be written as

$$R_s(\mathbf{W}, \mathbf{Z}, \Phi) = \left[ \min_{k \in \mathcal{K}} \{C_b(\mathbf{W}, \mathbf{Z}, \Phi) - C_{e,k}(\mathbf{W}, \mathbf{Z}, \Phi)\} \right]^+, \quad (5)$$

where  $[a]^+ \triangleq \max(a, 0)$ .

### B. Two CSI Error Scenarios

Since Bob is a registered and legitimate user, we assume that Bob's CSI is perfect while Eves' CSI is imperfect. Different from the communication system without IRS, there are two types of channels from Alice to Eves, which are the direct Alice-Eve channel  $\mathbf{h}_{ae,k}$  and the cascaded Alice-IRS-Eve channel  $\mathbf{G}_{ce,k}$ . Generally, these two kinds of channels are estimated separately and differently as shown in [18]. Thus, we consider two scenarios of CSI errors as follows.

1) *Scenario 1: Partial CSI Errors:* Due to the passive nature of IRS, the received signal from the reflecting link is weaker than that from the direct link, thus the probability of CSI errors on the cascaded Alice-IRS-Eve link is larger than the direct Alice-Eve link. Hence, in the first scenario, we assume the direct channel  $\mathbf{h}_{ae,k}$  is perfect, while the cascaded channel  $\mathbf{G}_{ce,k}$  is imperfect. The imperfect cascaded Alice-IRS-Eve channel can be presented as

$$\mathbf{G}_{ce,k} = \bar{\mathbf{G}}_{ce,k} + \Delta \mathbf{G}_{ce,k}, \quad \forall k \in \mathcal{K}, \quad (6)$$

where  $\bar{\mathbf{G}}_{ce,k}$  is the estimated value of  $\mathbf{G}_{ce,k}$  which is known to Alice, and  $\Delta \mathbf{G}_{ce,k}$  denotes the corresponding CSI error.

2) *Scenario 2: Full CSI Errors:* Since Eves are not legitimate users, we further consider full CSI errors on both the direct channels and cascaded channels for Eves in Scenario 2. In addition to the imperfect cascaded channel  $\mathbf{G}_{ce,k}$  in (6), the imperfect direct channel  $\mathbf{h}_{ae,k}$  can be described as

$$\mathbf{h}_{ae,k} = \bar{\mathbf{h}}_{ae,k} + \Delta \mathbf{h}_{ae,k}, \quad \forall k \in \mathcal{K}, \quad (7)$$

where  $\bar{\mathbf{h}}_{ae,k}$  is the estimated value of  $\mathbf{h}_{ae,k}$  which is known to Alice, and  $\Delta \mathbf{h}_{ae,k}$  denotes the corresponding CSI error.

We assume that the CSI errors originate from the channel estimation process, thus the statistical CSI error model is adopted here. Specifically, the CSI errors of  $\mathbf{h}_{ae,k}$  and  $\mathbf{G}_{ce,k}$  are assumed to be random and follow a CSCG distribution known a priori, i.e.,

$$\mathbf{g}_{he,k} \triangleq \Delta \mathbf{h}_{ae,k} \sim CN(\mathbf{0}, \Sigma_{he,k}), \quad \Sigma_{he,k} \geq \mathbf{0}, \quad (8a)$$

$$\mathbf{g}_{ge,k} \triangleq \text{vec}(\Delta \mathbf{G}_{ce,k}) \sim CN(\mathbf{0}, \Sigma_{ge,k}), \quad \Sigma_{ge,k} \geq \mathbf{0}, \quad (8b)$$

where  $\forall k \in \mathcal{K}$ , and  $\Sigma_{he,k} \in \mathbb{C}^{N_t \times N_t}$  and  $\Sigma_{ge,k} \in \mathbb{C}^{MN_t \times MN_t}$  are positive semidefinite covariance matrices of the CSI error. In addition,  $\mathbf{g}_{he,k}$  is independent of  $\mathbf{g}_{he,j}$  for any  $k \neq j$ , and  $\mathbf{g}_{ge,k}$  is independent of  $\mathbf{g}_{ge,j}$  for any  $k \neq j$ . The covariance matrices  $\Sigma_{he,k}$  and  $\Sigma_{ge,k}$  can describe the correlation property of CSI errors. For example, if these covariance matrices are identity matrices, we can regard that the CSI errors are uncorrelated.

### III. THE OCR TRANSMISSION DESIGN

To ensure the system security under Eves' CSI errors by sufficiently exploiting the spatial DoFs and AN, a robust transmission scheme should be designed. Thus, the transmit beamforming vector  $\mathbf{w}$ , AN spatial covariance  $\mathbf{Z}$  and IRS phase shifts  $\Phi$  are jointly optimized by an AO algorithm as follows, where the BTI, semi-definite relaxation (SDR) technique, and penalty CCP method [27] are leveraged.

#### A. Scenario 1: Partial CSI Errors

1) *Problem Formulation:* When the direct channel  $\mathbf{h}_{ae,k}$  of the Alice-Eve link is perfect, and the cascaded channel  $\Delta\mathbf{G}_{ce,k}$  of the Alice-IRS-Eve link is imperfect, we formulate an OCR-PM problem as

$$\min_{\mathbf{W}, \mathbf{Z}, \Phi} \text{Tr}(\mathbf{W} + \mathbf{Z}) \quad (9a)$$

$$\text{s.t. } C_b(\mathbf{W}, \mathbf{Z}, \Phi) \geq \log \gamma, \quad (9b)$$

$$\Pr_{\mathbf{g}_{ge,k}} \{C_{e,k}(\mathbf{W}, \mathbf{Z}, \Phi) \leq \log \beta\} \geq 1 - \rho_k, \forall k \in \mathcal{K}, \quad (9c)$$

$$\mathbf{Z} \geq 0, \quad (9d)$$

$$\mathbf{W} \geq 0, \quad (9e)$$

$$\text{rank}(\mathbf{W}) = 1, \quad (9f)$$

$$|\phi_m| = 1, m = 1, \dots, M, \quad (9g)$$

where  $\gamma \geq 1$  and  $\beta \geq 1$  are constant values, and  $\gamma \geq \beta$  is imposed to ensure a non-negative secrecy rate.  $\rho_k \in (0, 1]$  denotes the rate outage probability for the  $k$ th Eve. The chance constraint (9c) combined with (9b) guarantees that the probability that the secrecy rate is larger than  $\log \gamma - \log \beta$  is no less than  $1 - \rho_k$  in the presence of random CSI errors. By substituting (4a) and (4b) into (9), we can transform Problem (9) into

$$\min_{\mathbf{W}, \mathbf{Z}, \Phi} \text{Tr}(\mathbf{W} + \mathbf{Z}) \quad (10a)$$

$$\text{s.t. } \hat{\mathbf{h}}_b^H [\mathbf{W} - (\gamma - 1)\mathbf{Z}] \hat{\mathbf{h}}_b \geq (\gamma - 1)\sigma_b^2, \quad (10b)$$

$$\Pr_{\mathbf{g}_{ge,k}} \{ \hat{\mathbf{h}}_{e,k}^H [\mathbf{W} - (\beta - 1)\mathbf{Z}] \hat{\mathbf{h}}_{e,k} \leq \tilde{\sigma}_{e,k}^2 \} \geq 1 - \rho_k, \forall k \in \mathcal{K}, \quad (10c)$$

$$(9d), (9e), (9f), (9g). \quad (10d)$$

where  $\tilde{\sigma}_{e,k}^2 \triangleq (\beta - 1)\sigma_{e,k}^2$ . The main challenge for solving Problem (10) lies in the rate outage probability constraints (10c). To tackle it, we develop computable upper bounds for (10c) by using the BTI given in Lemma 1.

**Lemma 1.** (*Bernstein-Type Inequality* [34]) *For any  $(\mathbf{A}, \mathbf{u}, c) \in \mathbb{H}^n \times \mathbb{C}^n \times \mathbb{R}$ ,  $\mathbf{v} \sim \mathcal{CN}(\mathbf{0}, \mathbf{I}_n)$  and  $\rho \in (0, 1]$ , the following implication holds:*

$$\Pr_{\mathbf{v}} \{ \mathbf{v}^H \mathbf{A} \mathbf{v} + 2\text{Re}\{\mathbf{u}^H \mathbf{v}\} + c \geq 0 \} \geq 1 - \rho, \quad (11)$$

$$\iff \begin{cases} \text{Tr}(\mathbf{A}) - \sqrt{-2 \ln(\rho)} \cdot x + \ln(\rho) \cdot y + c \geq 0, \\ \left\| \begin{bmatrix} \text{vec}(\mathbf{A}) \\ \sqrt{2}\mathbf{u} \end{bmatrix} \right\|_2 \leq x, \\ \mathbf{y}\mathbf{I}_n + \mathbf{A} \geq \mathbf{0}, y \geq 0, \end{cases} \quad (12)$$

where  $x, y$  are the introduced slack variables.

As illustrated in [34], the (12) is a safe approximation of (11), and the tightness is achieved when the equalities in (12) hold.

2) *Reformulation:* To apply Lemma 1, we transform the probability constraints (10c) into the required form of (11) as follows. By substituting (3b), (6) into (10c), and defining  $\Xi_e \triangleq (\beta - 1)\mathbf{Z} - \mathbf{W}$ , the constraint on the probability of Eves' information leakage can be represented as

$$\begin{aligned} & \Pr_{\mathbf{g}_{ge,k}} \{ \hat{\mathbf{h}}_{e,k}^H [\mathbf{W} - (\beta - 1)\mathbf{Z}] \hat{\mathbf{h}}_{e,k} \leq \tilde{\sigma}_{e,k}^2 \} \\ &= \Pr_{\mathbf{g}_{ge,k}} \{ [\mathbf{h}_{ae,k}^H + \phi^T (\bar{\mathbf{G}}_{ce,k} + \Delta\mathbf{G}_{ce,k})] \Xi_e \\ & \quad \cdot [\mathbf{h}_{ae,k}^H + \phi^T (\bar{\mathbf{G}}_{ce,k} + \Delta\mathbf{G}_{ce,k})]^H + \tilde{\sigma}_{e,k}^2 \geq 0 \} \quad (13a) \\ &= \Pr_{\mathbf{g}_{ge,k}} \{ \underbrace{\phi^T \Delta\mathbf{G}_{ce,k} \Xi_e \Delta\mathbf{G}_{ce,k}^H \phi^*}_{f_{1,k}^{\text{Partial}}} \\ & \quad + 2\text{Re}[\underbrace{\phi^T \Delta\mathbf{G}_{ce,k} \Xi_e (\mathbf{h}_{ae,k} + \bar{\mathbf{G}}_{ce,k}^H \phi^*)}_{f_{2,k}^{\text{Partial}}}] \\ & \quad + \underbrace{(\mathbf{h}_{ae,k}^H + \phi^T \bar{\mathbf{G}}_{ce,k}) \Xi_e (\mathbf{h}_{ae,k} + \phi^T \bar{\mathbf{G}}_{ce,k})^H + \tilde{\sigma}_{e,k}^2}_{c_k^{\text{Partial}}} \geq 0 \}. \quad (13b) \end{aligned}$$

The matrix CSI error  $\Delta\mathbf{G}_{ce,k}$  in (13) can be substituted by its vectored form  $\mathbf{g}_{ge,k}$  defined in (8b), and then further substituted by its normalized and vectored form  $\mathbf{v}_{ge,k}$ . The  $\mathbf{v}_{ge,k}$  is related to  $\mathbf{g}_{ge,k}$  by  $\mathbf{g}_{ge,k} = \Sigma_{ge,k}^{1/2} \mathbf{v}_{ge,k}$ , where  $\mathbf{v}_{ge,k}$  is a standard Gaussian random vector, i.e.,  $\mathbf{v}_{ge,k} \sim \mathcal{CN}(\mathbf{0}, \mathbf{I}_{MN_t})$ , and  $\Sigma_{ge,k} = \Sigma_{ge,k}^{1/2} \Sigma_{ge,k}^{1/2}$ . Since  $\Sigma_{ge,k}$  is a positive semidefinite matrix, we have  $(\Sigma_{ge,k}^{1/2})^H = \Sigma_{ge,k}^{1/2}$  and  $(\Sigma_{ge,k}^{1/2})^T = (\Sigma_{ge,k}^{1/2})^*$ . Then, the  $f_{1,k}^{\text{Partial}}$  in (13b) can be reformulated into

$$\begin{aligned} f_{1,k}^{\text{Partial}} &= \text{Tr}(\Delta\mathbf{G}_{ce,k} \Xi_e \Delta\mathbf{G}_{ce,k}^H \phi^* \phi^T) = \text{Tr}(\Delta\mathbf{G}_{ce,k}^H \mathbf{E} \Delta\mathbf{G}_{ce,k} \Xi_e) \\ &\stackrel{(a)}{=} \text{vec}^H(\Delta\mathbf{G}_{ce,k}) (\Xi_e^T \otimes \mathbf{E}) \text{vec}(\Delta\mathbf{G}_{ce,k}) \triangleq \mathbf{v}_{ge,k}^H \mathbf{A}_{e,k} \mathbf{v}_{ge,k}, \quad (14) \end{aligned}$$

where  $\mathbf{E} \triangleq \phi^* \phi^T$ ,  $\mathbf{A}_{e,k} \triangleq \Sigma_{ge,k}^{1/2} (\Xi_e^T \otimes \mathbf{E}) \Sigma_{ge,k}^{1/2}$ , and (a) is obtained due to  $\text{Tr}(\mathbf{A}^H \mathbf{B} \mathbf{C} \mathbf{D}) = \text{vec}^H(\mathbf{A}) (\mathbf{D}^T \otimes \mathbf{B}) \text{vec}(\mathbf{C})$ . The expression of  $f_{2,k}^{\text{Partial}}$  in (13b) can be reformulated as

$$\begin{aligned} f_{2,k}^{\text{Partial}} &= \text{Tr}(\Delta\mathbf{G}_{ce,k} \Xi_e (\mathbf{h}_{ae,k} \phi^T + \bar{\mathbf{G}}_{ce,k}^H \mathbf{E})) \\ &\stackrel{(b)}{=} \text{vec}^H(\phi^* \mathbf{h}_{ae,k} + \mathbf{E} \bar{\mathbf{G}}_{ce,k}) (\Xi_e^T \otimes \mathbf{I}_M) \text{vec}(\Delta\mathbf{G}_{ce,k}) \\ &\triangleq \mathbf{u}_{e,k}^H \mathbf{v}_{ge,k}, \quad (15) \end{aligned}$$

where  $\mathbf{u}_{e,k} \triangleq \Sigma_{ge,k}^{1/2} (\Xi_e^* \otimes \mathbf{I}_M) \text{vec}(\phi^* \mathbf{h}_{ae,k} + \mathbf{E} \bar{\mathbf{G}}_{ce,k})$ , and (b) is obtained due to  $\text{Tr}(\mathbf{A} \mathbf{B} \mathbf{C}^H) = \text{vec}^H(\mathbf{C}) (\mathbf{B}^T \otimes \mathbf{I}) \text{vec}(\mathbf{A})$ .

By substituting (14) and (15) into (13b), the leakage data rate outage constraints (10c) for Eves become

$$(10c) \Leftrightarrow \Pr_{\mathbf{v}_{ge,k}} \left\{ \mathbf{v}_{ge,k}^H \mathbf{A}_{e,k} \mathbf{v}_{ge,k} + 2\text{Re}\{\mathbf{u}_{e,k}^H \mathbf{v}_{ge,k}\} + c_k^{\text{Partial}} \geq 0 \right\} \geq 1 - \rho_k, \forall k \in \mathcal{K}. \quad (16)$$

In (16), the outage probability w.r.t. the random CSI error  $\mathbf{g}_{ge,k}$  in (10c) is equivalently transformed into the outage probability of a real Gaussian quadratic form w.r.t.  $\mathbf{v}_{ge,k}$ , which facilitates the application of Lemma 1. It is also seen from (13) to (16) that Lemma 1 can be applied as long as the CSI error  $\mathbf{g}_{ge,k}$  follows the CSCG distribution and regardless of the correlation property of the CSI error. Then the chance constraint in (10c) can be conservatively approximated and

replaced by computable constraints according to Lemma 1, and Problem (10) becomes

$$\min_{\mathbf{W}, \mathbf{Z}, \phi, \mathbf{x}, \mathbf{y}} \text{Tr}(\mathbf{W} + \mathbf{Z}) \quad (17a)$$

$$\text{s.t. } \text{Tr}(\mathbf{A}_{e,k}) - \sqrt{-2 \ln(\rho_k)} \cdot x_k + \ln(\rho_k) \cdot y_k + c_k^{\text{Partial}} \geq 0, \quad (17b)$$

$$\left\| \begin{bmatrix} \text{vec}(\mathbf{A}_{e,k}) \\ \sqrt{2} \mathbf{u}_{e,k} \end{bmatrix} \right\|_2 \leq x_k, \quad (17c)$$

$$y_k \mathbf{I}_{MN_t} + \mathbf{A}_{e,k} \geq \mathbf{0}, y_k \geq 0, \quad (17d)$$

$$(10b), (9d), (9e), (9f), (9g), \quad (17e)$$

where  $\forall k \in \mathcal{K}$ ,  $\mathbf{x} = [x_1, x_2, \dots, x_K]^T$  and  $\mathbf{y} = [y_1, y_2, \dots, y_K]^T$  are introduced slack variables. We further simplify the  $\|\text{vec}(\mathbf{A}_{e,k})\|^2$  in constraint (17c) as

$$\begin{aligned} \|\text{vec}(\mathbf{A}_{e,k})\|^2 &= \|\mathbf{A}_{e,k}\|_F^2 = \text{Tr}[\mathbf{A}_{e,k} \mathbf{A}_{e,k}^H] \\ &= \text{Tr}[(\Xi_e^T \otimes \mathbf{E})^H \Sigma_{ge,k} (\Xi_e^T \otimes \mathbf{E}) \Sigma_{ge,k}] \\ &\stackrel{(c)}{=} \text{vec}^H(\Xi_e^T \otimes \mathbf{E}) (\Sigma_{ge,k}^T \otimes \Sigma_{ge,k}) \text{vec}(\Xi_e^T \otimes \mathbf{E}) \\ &= \text{vec}^H(\Xi_e^T \otimes \mathbf{E}) [(\Sigma_{ge,k}^{1/2T} \otimes \Sigma_{ge,k}^{1/2})^H (\Sigma_{ge,k}^{1/2T} \otimes \Sigma_{ge,k}^{1/2})] \text{vec}(\Xi_e^T \otimes \mathbf{E}) \\ &= \left\| (\Sigma_{ge,k}^{1/2T} \otimes \Sigma_{ge,k}^{1/2}) \text{vec}(\Xi_e^T \otimes \mathbf{E}) \right\|^2, \end{aligned} \quad (18)$$

where (c) is obtained by invoking the identity  $\text{Tr}(\mathbf{A}^H \mathbf{B} \mathbf{C} \mathbf{D}) = \text{vec}^H(\mathbf{A}) (\mathbf{D}^T \otimes \mathbf{B}) \text{vec}(\mathbf{C})$ .

By substituting the expressions of (18) into (17c), we have the OCR-PM Problem in (17) written more explicitly as

$$\min_{\mathbf{W}, \mathbf{Z}, \phi, \mathbf{x}, \mathbf{y}} \text{Tr}(\mathbf{W} + \mathbf{Z}) \quad (19a)$$

$$\text{s.t. } \text{Tr}(\Sigma_{ge,k}^{1/2} (\Xi_e^T \otimes \mathbf{E}) \Sigma_{ge,k}^{1/2}) - \sqrt{-2 \ln(\rho_k)} \cdot x_k + \ln(\rho_k) \cdot y_k + c_k^{\text{Partial}} \geq 0, \quad (19b)$$

$$\left\| \begin{bmatrix} (\Sigma_{ge,k}^{1/2T} \otimes \Sigma_{ge,k}^{1/2}) \text{vec}(\Xi_e^T \otimes \mathbf{E}) \\ \sqrt{2} \Sigma_{ge,k}^{1/2} (\Xi_e^* \otimes \mathbf{I}_M) \text{vec}(\phi^* \mathbf{h}_{ae,k}^H + \mathbf{E} \tilde{\mathbf{G}}_{ce,k}) \end{bmatrix} \right\|_2 \leq x_k, \quad (19c)$$

$$y_k \mathbf{I}_{MN_t} + \Sigma_{ge,k}^{1/2} (\Xi_e^T \otimes \mathbf{E}) \Sigma_{ge,k}^{1/2} \geq \mathbf{0}, y_k \geq 0, \quad (19d)$$

$$(10b), (9d), (9e), (9f), (9g), \quad (19e)$$

where  $\forall k \in \mathcal{K}$ . The chance constraints are removed in Problem (19). However, the optimization variables  $\{\mathbf{W}, \mathbf{Z}\}$  and  $\phi$  are coupled. We propose to tackle the coupling via the AO method. The variables  $\{\mathbf{W}, \mathbf{Z}\}$  and  $\phi$  are updated alternately.

3) *Optimization of Transmit Beamforming and AN*: Obviously, when  $\phi$  is fixed, the  $c_k^{\text{Partial}}$ ,  $\text{Tr}(\mathbf{A}_k)$ ,  $\|\text{vec}(\mathbf{A}_k)\|^2$  and  $\|\mathbf{u}_{ce,k}\|$  are all convex functions of  $\Xi_e$ . Then, Problem (19) cannot be solved efficiently only due to the nonconvexity of  $\text{rank}(\mathbf{W}) = 1$  in (9f). By removing (9f), the  $\{\mathbf{W}, \mathbf{Z}\}$  can be solved by the SDR, and the corresponding optimization problem is

$$\min_{\mathbf{W}, \mathbf{Z}, \mathbf{x}, \mathbf{y}} \text{Tr}(\mathbf{W} + \mathbf{Z}) \quad (20a)$$

$$\text{s.t. } (19b), (19c), (19d), (10b), (9d), (9e), \forall k \in \mathcal{K}. \quad (20b)$$

The SDR in Problem (20) is tight, which means that the solved  $\mathbf{W}$  always satisfies  $\text{rank}(\mathbf{W}) = 1$ . Then the beamforming vector  $\mathbf{w}$  can be recovered from  $\mathbf{W} = \mathbf{w} \mathbf{w}^H$  by performing the Cholesky decomposition. The tightness of the SDR is proved in the following Theorem.

**Theorem 1.** Assume that the optimal solution of Problem (20) is  $(\mathbf{W}^*, \mathbf{Z}^*, \mathbf{x}^*, \mathbf{y}^*)$ , where  $\text{rank}(\mathbf{W}^*) \geq 1$ . Then there always exists another optimal solution of Problem (20), denoted as  $(\tilde{\mathbf{W}}^*, \tilde{\mathbf{Z}}^*, \tilde{\mathbf{x}}^*, \tilde{\mathbf{y}}^*)$ , which not only has the same objective value as  $(\mathbf{W}^*, \mathbf{Z}^*, \mathbf{x}^*, \mathbf{y}^*)$ , but also satisfies the rank-one constraint, i.e.,  $\text{rank}(\tilde{\mathbf{W}}^*) = 1$ .

*Proof.* The proving process is the same as that in Appendix A of [33] after replacing the channel vector  $\hat{\mathbf{h}}_b$  in [33] by  $\hat{\mathbf{h}}_b^H \triangleq \mathbf{h}_{ab}^H + \phi^T \mathbf{G}_{cb}$  defined in (3a) of this paper.  $\square$

4) *Optimization of Phase Shifts at the IRS*: For given  $\mathbf{W}$  and  $\mathbf{Z}$ , the objective function of Problem (19) is irrelevant with  $\phi$ . To achieve better convergence, slack variables are introduced, and the data rate constraints in (10b) and in (13a) can be modified respectively as

$$\hat{\mathbf{h}}_b^H [\mathbf{W} - (\gamma - 1) \mathbf{Z}] \hat{\mathbf{h}}_b \geq \bar{\sigma}_b^2 + \delta_0, \delta_0 \geq 0, \quad (21a)$$

$$\begin{aligned} [\mathbf{h}_{ae,k}^H + \phi^T (\tilde{\mathbf{G}}_{ce,k} + \Delta \mathbf{G}_{ce,k})] \Xi_e [\mathbf{h}_{ae,k}^H + \phi^T (\tilde{\mathbf{G}}_{ce,k} + \Delta \mathbf{G}_{ce,k})]^H \\ + \bar{\sigma}_{e,k}^2 - \delta_k \geq 0, \delta_k \geq 0, \forall k \in \mathcal{K}, \end{aligned} \quad (21b)$$

where  $\bar{\sigma}_b^2 \triangleq (\gamma - 1) \sigma_b^2$ , and  $\delta_0, \delta_k, \forall k \in \mathcal{K}$  are slack variables.

Then, the outage probability of (21b) can be safely approximated again by invoking Lemma 1, and the optimization problem for  $\phi$  can be written as

$$\max_{\phi, \delta, \mathbf{x}, \mathbf{y}} \sum_{k=0}^K \delta_k \quad (22a)$$

$$\text{s.t. } \text{Tr}(\Sigma_{ge,k}^{1/2} (\Xi_e^T \otimes \mathbf{E}) \Sigma_{ge,k}^{1/2}) - \sqrt{-2 \ln(\rho_k)} \cdot x_k + \ln(\rho_k) \cdot y_k + c_k^{\text{Partial}} - \delta_k \geq 0, \forall k \in \mathcal{K}, \quad (22b)$$

$$(19c), (19d), (9g), \forall k \in \mathcal{K}, \quad (22c)$$

$$\hat{\mathbf{h}}_b^H [\mathbf{W} - (\gamma - 1) \mathbf{Z}] \hat{\mathbf{h}}_b \geq \bar{\sigma}_b^2 + \delta_0, \quad (22d)$$

$$\delta \geq \mathbf{0}, \quad (22e)$$

where  $\delta = [\delta_0, \delta_1, \delta_2, \dots, \delta_K]^T$  are slack variables.

It is noted that Problem (22) cannot be transformed into the form with optimization variable of  $\mathbf{E} = \phi^* \phi^T$ , thus cannot be solved by SDR as in [33] where the direct communication links are obstructed. Due to the existence of direct links in this paper, the optimization variable can only be  $\phi$  instead of  $\mathbf{E}$ , then almost all constraints in Problem (22) become nonconvex. To tackle this challenge, we try to transform Problem (22) into a convex problem w.r.t.  $\phi$ , and propose novel mathematical manipulations on these constraints in four steps as follows.

**Step 1:** Transform the constraint of (22b) into convex ones.

(1) We note that the  $c_k^{\text{Partial}}$  in (22b) is non-concave w.r.t.  $\phi$  due to the fact that  $\tilde{\mathbf{G}}_{ce,k} \Xi_e \tilde{\mathbf{G}}_{ce,k}^H$  is a non-negative semidefinite matrix. To address this issue, the term  $c_k^{\text{Partial}}$  can be equivalently rewritten as

$$\begin{aligned} c_k^{\text{Partial}}(\phi) &= \phi^T (\tilde{\mathbf{G}}_{ce,k} \Xi_e \tilde{\mathbf{G}}_{ce,k}^H - c_{c,k} \mathbf{I}_M) \phi^* \\ &+ 2\text{Re}\{\phi^T \tilde{\mathbf{G}}_{ce,k} \Xi_e \mathbf{h}_{ae,k}\} + \mathbf{h}_{ae,k}^H \Xi_e \mathbf{h}_{ae,k} + M c_{c,k} + \bar{\sigma}_{e,k}^2, \end{aligned} \quad (23)$$

where the unit-modulus property  $\phi^T \phi^* = M$  is utilized, and we set  $c_{c,k} = \lambda_{\max}\{\tilde{\mathbf{G}}_{ce,k} (\beta - 1) \mathbf{Z} \tilde{\mathbf{G}}_{ce,k}^H\}$ . Since  $\Xi_e = (\beta - 1) \mathbf{Z} - \mathbf{W}$ , we have  $\tilde{\mathbf{G}}_{ce,k} \Xi_e \tilde{\mathbf{G}}_{ce,k}^H - c_{c,k} \mathbf{I}_M \leq \mathbf{0}$ , and the expression of  $c_k^{\text{Partial}}(\phi)$  in (23) becomes concave w.r.t.  $\phi$ .

(2) The term  $\text{Tr}(\Sigma_{ge,k}^{1/2}(\Xi_e^T \otimes \mathbf{E})\Sigma_{ge,k}^{1/2})$  in (22b) is equivalently transformed as

$$\begin{aligned} & \text{Tr}(\Sigma_{ge,k}^{1/2}(\Xi_e^T \otimes \mathbf{E})\Sigma_{ge,k}^{1/2}) \\ & \stackrel{(d)}{=} \text{vec}^H(\Sigma_{ge,k}^{1/2})(\mathbf{I}_{MN_t} \otimes (\Xi_e^T \otimes (\phi^* \phi^T)))\text{vec}(\Sigma_{ge,k}^{1/2}) \\ & = \text{vec}^H(\Sigma_{gx,k})(\mathbf{I}_{MN_t} \otimes \Xi_e^T \otimes (\phi^* \phi^T))\text{vec}(\Sigma_{gx,k}) \\ & \stackrel{(e)}{=} \phi^T \Sigma_{gx,k} (\mathbf{I}_{MN_t} \otimes \Xi_e) \Sigma_{gx,k}^H \phi^*, \end{aligned} \quad (24)$$

where the (d) is obtained by invoking  $\text{Tr}(\mathbf{A}^H \mathbf{B} \mathbf{C} \mathbf{D}) = \text{vec}^H(\mathbf{A})(\mathbf{D}^T \otimes \mathbf{B})\text{vec}(\mathbf{C})$ . The (e) follows from the equality  $\text{vec}(\mathbf{X})^H (\mathbf{B}^T \otimes \mathbf{c} \mathbf{a}^H) \text{vec}(\mathbf{X}) = \mathbf{a}^H \mathbf{X} \mathbf{B} \mathbf{X}^H \mathbf{c}$ , where  $\mathbf{B}^T = \mathbf{I}_{MN_t} \otimes \Xi_e^T$ ,  $\mathbf{c} = \phi^*$ ,  $\mathbf{a}^H = \phi^T$ , and  $\mathbf{X} = \Sigma_{gx,k}$ . When using this identity, the matrix dimension must be matched. Specifically,  $\Sigma_{gx,k} \in \mathbb{C}^{M \times MN_t^2}$  is obtained by remapping the column vector  $\text{vec}(\Sigma_{ge,k}^{1/2}) \in \mathbb{C}^{(MN_t)^2 \times 1}$  into an  $M \times MN_t^2$  dimensional matrix, which can be expressed as  $\Sigma_{gx,k} = \text{unvec}[\text{vec}(\Sigma_{ge,k}^{1/2})]_{M \times MN_t^2}$ .

It is noted that although (24) is a quadratic form w.r.t.  $\phi$ , but it is still non-concave w.r.t.  $\phi$ . Similarly, we reformulate (24) into a concave form as

$$\begin{aligned} & \text{Tr}(\Sigma_{ge,k}^{1/2}(\Xi_e^T \otimes \mathbf{E})\Sigma_{ge,k}^{1/2}) \\ & = \phi^T [\Sigma_{gx,k} (\mathbf{I}_{MN_t} \otimes \Xi_e) \Sigma_{gx,k}^H - c_{t,k} \mathbf{I}_M] \phi^* + M c_{t,k}, \end{aligned} \quad (25)$$

where  $c_{t,k} = \lambda_{\max}\{\Sigma_{gx,k} (\mathbf{I}_{MN_t} \otimes ((\beta - 1)\mathbf{Z})) \Sigma_{gx,k}^H\}$ .

(3) By substituting (23) and (25) into (22b), we have the following convex constraint:

$$(22b) \Leftrightarrow c_{ic,k}^{\text{partial}}(\phi) \geq 0, \quad (26)$$

where

$$\begin{aligned} c_{ic,k}^{\text{partial}}(\phi) & \triangleq \phi^T [\Sigma_{gx,k} (\mathbf{I}_{MN_t} \otimes \Xi_e) \Sigma_{gx,k}^H - c_{t,k} \mathbf{I}_M] \phi^* + M c_{t,k} \\ & - \sqrt{-2 \ln(\rho_k)} \cdot x_k + \ln(\rho_k) \cdot y_k + \phi^T (\bar{\mathbf{G}}_{ce,k} \Xi_e \bar{\mathbf{G}}_{ce,k}^H - c_{c,k} \mathbf{I}_M) \phi^* \\ & + 2 \text{Re}\{\phi^T \bar{\mathbf{G}}_{ce,k} \Xi_e \mathbf{h}_{ae,k}\} + \mathbf{h}_{ae,k}^H \Xi_e \mathbf{h}_{ae,k} + M c_{c,k} + \bar{\sigma}_{e,k}^2 - \delta_k. \end{aligned} \quad (27)$$

As a result, the non-convex constraint in (22b) is transformed into a convex one in (26).

**Step 2:** Transform the constraint in (19c) to be convex.

(1) In constraint (19c), we can find an upper bound of  $(\Sigma_{ge,k}^{1/2T} \otimes \Sigma_{ge,k}^{1/2})\text{vec}(\Xi_e^T \otimes \mathbf{E})$  as

$$\begin{aligned} & \left\| (\Sigma_{ge,k}^{1/2T} \otimes \Sigma_{ge,k}^{1/2})\text{vec}(\Xi_e^T \otimes \mathbf{E}) \right\| \\ & \stackrel{(f)}{\leq} \left\| (\Sigma_{ge,k}^{1/2T} \otimes \Sigma_{ge,k}^{1/2}) \right\|_2 \left\| \text{vec}(\Xi_e^T \otimes \mathbf{E}) \right\|_2, \end{aligned} \quad (28a)$$

$$\stackrel{(g)}{=} \lambda_{\max}(\Sigma_{ge,k}^{1/2T} \otimes \Sigma_{ge,k}^{1/2}) M \|\Xi_e\|_F, \quad (28b)$$

where the (f) is obtained due to  $\|\mathbf{A}\mathbf{x}\|_2 \leq \|\mathbf{A}\|_2 \|\mathbf{x}\|_2$ , and  $\|\mathbf{A}\|_2 = \sigma_{\max}\{\mathbf{A}\}$  denotes the spectral norm of matrix  $\mathbf{A}$ . The equality  $\|\mathbf{A}\mathbf{x}\|_2 = \|\mathbf{A}\|_2 \|\mathbf{x}\|_2$  holds when the matrix  $\mathbf{A}$  is an unitary matrix. The step (g) is obtained due to  $\sigma_{\max}\{\mathbf{A}\} = \lambda_{\max}\{\mathbf{A}\}$  if  $\mathbf{A} \in \mathbb{H}^n$ , and the following property:

$$\begin{aligned} \left\| \text{vec}(\Xi_e^T \otimes \mathbf{E}) \right\|^2 & = \text{Tr}[(\Xi_e^T \otimes \mathbf{E})(\Xi_e^* \otimes \mathbf{E})] = \text{Tr}[(\Xi_e^T \Xi_e^*) \otimes (\mathbf{E}\mathbf{E})] \\ & = \text{Tr}[(\Xi_e^T \Xi_e^*)] \text{Tr}[(\mathbf{E}\mathbf{E})] = M^2 \|\Xi_e\|_F^2. \end{aligned} \quad (29)$$

It is observed that when  $\Sigma_{ge,k}^{1/2T} \otimes \Sigma_{ge,k}^{1/2}$  is an unitary matrix, e.g.,  $\Sigma_{ge,k} = \mathbf{I}_{MN_t}$ , the inequality of (f) in (28a) holds with equality.

(2) In (19c), we can find an upper bound of  $\Sigma_{ge,k}^{1/2}(\Xi_e^* \otimes \mathbf{I}_M)\text{vec}(\phi^* \mathbf{h}_{ae,k}^H + \mathbf{E} \bar{\mathbf{G}}_{ce,k})$  as

$$\begin{aligned} & \left\| \Sigma_{ge,k}^{1/2}(\Xi_e^* \otimes \mathbf{I}_M)\text{vec}[\phi^* (\mathbf{h}_{ae,k}^H + \phi^T \bar{\mathbf{G}}_{ce,k})] \right\| \\ & \stackrel{(h)}{\leq} \lambda_{\max}(\Sigma_{ge,k}^{1/2}(\Xi_e^* \otimes \mathbf{I}_M)) \left\| (\mathbf{h}_{ae,k}^* + \bar{\mathbf{G}}_{ce,k}^T \phi) \otimes \phi^* \right\|, \end{aligned} \quad (30a)$$

$$\stackrel{(i)}{=} \lambda_{\max}(\Sigma_{ge,k}^{1/2}(\Xi_e^* \otimes \mathbf{I}_M)) \sqrt{M} \left\| (\mathbf{h}_{ae,k}^* + \bar{\mathbf{G}}_{ce,k}^T \phi) \right\|, \quad (30b)$$

where the (h) is obtained due to  $\|\mathbf{A}\mathbf{x}\|_2 \leq \|\mathbf{A}\|_2 \|\mathbf{x}\|_2$  and  $\text{vec}(\mathbf{a}\mathbf{b}^T) = \mathbf{b} \otimes \mathbf{a}$ . It is observed that when  $\Sigma_{ge,k}^{1/2}(\Xi_e^* \otimes \mathbf{I}_M)$  is an unitary matrix, the inequality of (h) in (30a) becomes an equality. The (i) is obtained by using Lemma 2 below.

**Lemma 2.** For any two vectors  $\mathbf{b}$  and  $\mathbf{c}$ , we have  $\|\mathbf{b} \otimes \mathbf{c}\| = \|\mathbf{b}\| \|\mathbf{c}\|$ . In particular, we have  $\|\mathbf{a} \otimes \mathbf{a}\| = \mathbf{a}^H \mathbf{a}$ .

**Proof.** Please refer to Appendix A.  $\square$

Thus, by utilizing the upper bounds in (28b) and (30b), the constraint in (19c) can be approximated by

$$\left\| \frac{\lambda_{\max}(\Sigma_{ge,k}^{1/2T} \otimes \Sigma_{ge,k}^{1/2}) M \text{vec}(\Xi_e)}{\sqrt{2M} \lambda_{\max}(\Sigma_{ge,k}^{1/2}(\Xi_e^* \otimes \mathbf{I}_M)) (\mathbf{h}_{ae,k}^* + \bar{\mathbf{G}}_{ce,k}^T \phi)} \right\| \leq x_k. \quad (31)$$

Obviously, the constraint (31) is convex, and we have (19c)  $\Leftrightarrow$  (31). When both  $(\Sigma_{ge,k}^{1/2T} \otimes \Sigma_{ge,k}^{1/2})$  and  $\Sigma_{ge,k}^{1/2}(\Xi_e^* \otimes \mathbf{I}_M)$  are unitary matrices, we have (19c)  $\Leftrightarrow$  (31).

**Step 3:** Transform the constraint of (19d) to be convex.

Since  $\beta \geq 1$ , we note that  $c_g = \lambda_{\max}\{(\beta - 1)\mathbf{Z}\} \geq 0$ . By substituting  $\Xi_e = (\beta - 1)\mathbf{Z} - \mathbf{W}$  into (19d), and adding  $\Sigma_{ge,k}^{1/2}((c_g \mathbf{I}_{N_t})^T \otimes \mathbf{E})\Sigma_{ge,k}^{1/2}$  on both sides of (19d), we have

$$\begin{aligned} y_k \mathbf{I}_{N_t M} + \Sigma_{ge,k}^{1/2}((c_g \mathbf{I}_{N_t})^T \otimes \mathbf{E})\Sigma_{ge,k}^{1/2} & \geq \Sigma_{ge,k}^{1/2}(\mathbf{W}^T \otimes \mathbf{E})\Sigma_{ge,k}^{1/2} \\ & + \Sigma_{ge,k}^{1/2}[(c_g \mathbf{I}_{N_t} - (\beta - 1)\mathbf{Z})^T \otimes \mathbf{E}]\Sigma_{ge,k}^{1/2}, \end{aligned} \quad (32a)$$

$$\begin{aligned} & \Leftrightarrow y_k \mathbf{I}_{N_t M} + \Sigma_{ge,k}^{1/2}((c_g \mathbf{I}_{N_t})^T \otimes \mathbf{E})\Sigma_{ge,k}^{1/2} \\ & \geq \Sigma_{ge,k}^{1/2}[(\mathbf{W} + \mathbf{Z}_g)^T \otimes \mathbf{E}]\Sigma_{ge,k}^{1/2}, \end{aligned} \quad (32b)$$

$$\begin{aligned} & \Leftrightarrow y_k + \lambda_{\min}\{\Sigma_{ge,k}^{1/2}((c_g \mathbf{I}_{N_t})^T \otimes \mathbf{E})\Sigma_{ge,k}^{1/2}\} \\ & \geq \lambda_{\max}\{\Sigma_{ge,k}^{1/2}(\mathbf{V}_{wgz}^* \otimes \phi^*)(\mathbf{V}_{wgz}^T \otimes \phi^T)\Sigma_{ge,k}^{1/2}\}, \end{aligned} \quad (32c)$$

$$\Leftrightarrow y_k \geq \left\| \Sigma_{ge,k}^{1/2}(\mathbf{V}_{wgz}^* \otimes \phi^*) \right\|_2^2, \quad (32d)$$

where  $\mathbf{Z}_g \triangleq c_g \mathbf{I}_{N_t} - (\beta - 1)\mathbf{Z}$ ,  $\mathbf{W} + \mathbf{Z}_g \triangleq \mathbf{V}_{wgz} \mathbf{V}_{wgz}^H$ ,  $\mathbf{Z}_g \geq \mathbf{0}$  and  $\mathbf{W} + \mathbf{Z}_g \geq \mathbf{0}$ . The step in (j) is obtained because  $\lambda_{\min}\{\Sigma_{ge,k}^{1/2}((c_g \mathbf{I}_{N_t})^T \otimes \mathbf{E})\Sigma_{ge,k}^{1/2}\} = 0$ , and the proof is given in Appendix B. The constraint (32d) is convex, and can be equivalently written as

$$(19d) \Leftrightarrow \left\| \Sigma_{ge,k}^{1/2}(\mathbf{V}_{wgz}^* \otimes \phi^*) \right\|_2 - \sqrt{y_k} \leq 0, \quad (33)$$

which is convex.

**Step 4:** Transform the constraint of (22d) to be convex.

By substituting (3a) into the data rate requirement for Bob in (22d), and defining  $\Xi_b \triangleq (\gamma - 1)\mathbf{Z} - \mathbf{W}$ , the constraint (22d) can be recast as

$$(22d) \Leftrightarrow -(\mathbf{h}_{ab}^H + \phi^T \mathbf{G}_{cb}) \Xi_b (\mathbf{h}_{ab} + \mathbf{G}_{cb}^H \phi^*) \geq \bar{\sigma}_b^2 + \delta_0, \quad (34a)$$

$$\begin{aligned} & \Leftrightarrow \phi^T \mathbf{G}_{cb} \Xi_b \mathbf{G}_{cb}^H \phi^* + 2 \text{Re}\{\phi^T \mathbf{G}_{cb} \Xi_b \mathbf{h}_{ab}\} \\ & + \mathbf{h}_{ab}^H \Xi_b \mathbf{h}_{ab} + \bar{\sigma}_b^2 + \delta_0 \leq 0. \end{aligned} \quad (34b)$$

To build the convexity of (34b), it is reformulated as

$$(22d) \Leftrightarrow b(\phi) \leq 0, \quad (35)$$

where  $c_b = \lambda_{\min}\{\mathbf{G}_{cb}\mathbf{\Xi}_b\mathbf{G}_{cb}^H\}$  and

$$b(\phi) \triangleq \phi^T [\mathbf{G}_{cb}\mathbf{\Xi}_b\mathbf{G}_{cb}^H - c_b\mathbf{I}_M] \phi^* + 2\text{Re}\{\phi^T \mathbf{G}_{cb}\mathbf{\Xi}_b \mathbf{h}_{ab}\} + \mathbf{h}_{ab}^H \mathbf{\Xi}_b \mathbf{h}_{ab} + \tilde{\sigma}_b^2 + \delta_0 + M c_b. \quad (36)$$

Finally, based on the mathematical manipulations of **Steps 1-4** above, Problem (22) is reformulated into

$$\max_{\phi, \delta, \mathbf{x}, \mathbf{y}} \sum_{k=0}^K \delta_k \quad (37a)$$

$$\text{s.t. (26), (31), (33), (35), (22e), } \forall k \in \mathcal{K}, \quad (37b)$$

$$|\phi_m| = 1, m = 1, \dots, M, \quad (37c)$$

which is nonconvex only due to the unit-modulus constraint in (37c). To tackle it, the penalty CCP algorithm [27] is leveraged. By minimizing the sum of non-negative slack variables in  $\mathbf{b} = [b_1, \dots, b_{2M}]^T$ , the constraint of (37c) is equivalent to  $|\phi_m|^2 \leq 1 + b_{M+m}$  and  $|\phi_m|^2 \geq 1 - b_m$ , where the former is convex and the latter is nonconvex. By using the first-order Taylor expansion, the nonconvex constraint  $|\phi_m|^2 \geq 1 - b_m$  is transformed into an affine constraint  $|\phi_m^{(t)}|^2 - 2\text{Re}\{\phi_m^* \phi_m^{(t)}\} \leq b_m - 1$ , where  $\phi_m^{(t)}$  is the phase shift value obtained in last iteration ( $t$ th iteration). Then, Problem (37) can be transformed into an iterative optimization process, and the optimization problem at the ( $t+1$ )th iteration is

$$\max_{\phi, \delta, \mathbf{b}, \mathbf{x}, \mathbf{y}} \sum_{k=0}^K \delta_k - \lambda^{(t+1)} \sum_{l=1}^{2M} b_l \quad (38a)$$

$$\text{s.t. (26), (31), (33), (35), (22e), } \forall k \in \mathcal{K}, \quad (38b)$$

$$|\phi_m^{(t)}|^2 - 2\text{Re}\{\phi_m^* \phi_m^{(t)}\} \leq b_m - 1, m = 1, \dots, M, \quad (38c)$$

$$|\phi_m|^2 \leq 1 + b_{M+m}, m = 1, \dots, M, \quad (38d)$$

$$\mathbf{b} \geq \mathbf{0}, \quad (38e)$$

where  $\sum_{l=1}^{2M} b_l$  is added into the objective function as a penalty term, and  $\lambda^{(t+1)}$  is employed as the regularization factor to control the feasibility of the constraints.

## B. Scenario 2: Full CSI Errors

1) *Problem Reformulation:* Considering the full statistical CSI error model in (8), the outage probability constraints in (10c) for Eves' leaked data rate can be extended to

$$\Pr_{\mathbf{g}_{he,k}, \mathbf{g}_{ge,k}} \left\{ \hat{\mathbf{h}}_{e,k}^H [\mathbf{W} - (\beta - 1)\mathbf{Z}] \hat{\mathbf{h}}_{e,k} \leq (\beta - 1)\sigma_{e,k}^2 \right\} \geq 1 - \rho_k, \forall k \in \mathcal{K}. \quad (39)$$

To apply Lemma 1, we transform the probability constraints (39) into the required form of (11) as follows. By substituting

(3b), (6) and (7) into (39), the (39) can be reformulated into

$$\begin{aligned} & \Pr_{\mathbf{g}_{he,k}, \mathbf{g}_{ge,k}} \left\{ [\mathbf{h}_{ae,k}^H + \Delta \mathbf{h}_{ae,k}^H + \phi^T (\tilde{\mathbf{G}}_{ce,k} + \Delta \mathbf{G}_{ce,k})] \mathbf{\Xi}_e \right. \\ & \cdot [\mathbf{h}_{ae,k}^H + \Delta \mathbf{h}_{ae,k}^H + \phi^T (\tilde{\mathbf{G}}_{ce,k} + \Delta \mathbf{G}_{ce,k})]^H + \tilde{\sigma}_{e,k}^2 \geq 0 \left. \right\} \quad (40a) \\ & = \Pr_{\mathbf{g}_{he,k}, \mathbf{g}_{ge,k}} \left\{ \underbrace{(\Delta \mathbf{h}_{ae,k}^H + \phi^T \Delta \mathbf{G}_{ce,k}) \mathbf{\Xi}_e (\Delta \mathbf{h}_{ae,k} + \Delta \mathbf{G}_{ce,k}^H \phi^*)}_{f_{1,k}^{\text{full}}} \right. \\ & \quad \left. + 2\text{Re}\{(\mathbf{h}_{ae,k}^H + \phi^T \tilde{\mathbf{G}}_{ce,k}) \mathbf{\Xi}_e (\Delta \mathbf{h}_{ae,k} + \Delta \mathbf{G}_{ce,k}^H \phi^*)\} \right. \\ & \quad \left. + \underbrace{(\mathbf{h}_{ae,k}^H + \phi^T \tilde{\mathbf{G}}_{ce,k}) \mathbf{\Xi}_e (\mathbf{h}_{ae,k}^H + \phi^T \tilde{\mathbf{G}}_{ce,k})^H + \tilde{\sigma}_{e,k}^2}_{f_{2,k}^{\text{full}}} \geq 0 \right\}. \quad (40b) \end{aligned}$$

Similar as the manipulations under (13), the CSI error matrix  $\Delta \mathbf{G}_{ce,k}$  in (40) can be substituted by its normalized and vectored form  $\mathbf{v}_{ge,k}$ , and the CSI error vector  $\Delta \mathbf{h}_{ae,k}$  in (40) can also be substituted by its normalized form  $\mathbf{v}_{he,k}$ . According to (8a), the  $\mathbf{v}_{he,k}$  is related to  $\mathbf{g}_{he,k}$  by  $\mathbf{g}_{he,k} \triangleq \Delta \mathbf{h}_{ae,k} = \Sigma_{he,k}^{1/2} \mathbf{v}_{he,k}$ , where  $\mathbf{v}_{he,k} \sim \mathcal{CN}(\mathbf{0}, \mathbf{I}_{N_t})$  and  $\Sigma_{he,k} = \Sigma_{he,k}^{1/2} \Sigma_{he,k}^{1/2}$ . Since  $\Sigma_{he,k}$  is a semidefinite matrix, we have  $(\Sigma_{he,k}^{1/2})^H = \Sigma_{he,k}^{1/2}$  and  $(\Sigma_{he,k}^{1/2})^T = (\Sigma_{he,k}^{1/2})^*$ . By invoking the identity  $\text{Tr}(\mathbf{A}\mathbf{B}\mathbf{C}\mathbf{D}) = \text{vec}^T(\mathbf{D})(\mathbf{A} \otimes \mathbf{C}^T) \text{vec}(\mathbf{B}^T)$ , the expression of  $f_{2,k}^{\text{full}}$  in (40b) can be rewritten as

$$\begin{aligned} f_{2,k}^{\text{full}} &= (\mathbf{h}_{ae,k}^H + \phi^T \tilde{\mathbf{G}}_{ce,k}) \mathbf{\Xi}_e \mathbf{g}_{he,k} + \text{Tr}\{[\mathbf{\Xi}_e \Delta \mathbf{G}_{ce,k}^H \phi^* (\mathbf{h}_{ae,k}^H + \phi^T \tilde{\mathbf{G}}_{ce,k})]\} \\ &= (\mathbf{h}_{ae,k}^H + \phi^T \tilde{\mathbf{G}}_{ce,k}) \mathbf{\Xi}_e \mathbf{g}_{he,k} \\ & \quad + \text{vec}^T(\mathbf{h}_{ae,k}^H + \phi^T \tilde{\mathbf{G}}_{ce,k}) (\mathbf{\Xi}_e \otimes \phi^H) \text{vec}(\Delta \mathbf{G}_{ce,k}^*) \\ &= (\mathbf{h}_{ae,k}^H + \phi^T \tilde{\mathbf{G}}_{ce,k}) \mathbf{\Xi}_e \Sigma_{he,k}^{1/2} \mathbf{v}_{he,k} \\ & \quad + (\mathbf{h}_{ae,k}^H + \phi^T \tilde{\mathbf{G}}_{ce,k}) (\mathbf{\Xi}_e \otimes \phi^H) \Sigma_{ge,k}^{1/2*} \mathbf{v}_{ge,k}^* \\ & \triangleq \tilde{\mathbf{u}}_{e,k}^H \tilde{\mathbf{v}}_{e,k}, \quad (41) \end{aligned}$$

where  $\tilde{\mathbf{v}}_{e,k} = [\mathbf{v}_{he,k}^H, \mathbf{v}_{ge,k}^T]^H$ , and

$$\begin{aligned} \tilde{\mathbf{u}}_{e,k} &= \begin{bmatrix} \Sigma_{he,k}^{1/2} \mathbf{\Xi}_e (\mathbf{h}_{ae,k} + \tilde{\mathbf{G}}_{ce,k}^H \phi^*) \\ \Sigma_{ge,k}^{1/2T} (\mathbf{\Xi}_e \otimes \phi) (\mathbf{h}_{ae,k} + \tilde{\mathbf{G}}_{ce,k}^H \phi^*) \end{bmatrix} \\ &\stackrel{(k)}{=} \begin{bmatrix} \Sigma_{he,k}^{1/2} \mathbf{\Xi}_e (\mathbf{h}_{ae,k} + \tilde{\mathbf{G}}_{ce,k}^H \phi^*) \\ \Sigma_{ge,k}^{1/2T} ([\mathbf{\Xi}_e (\mathbf{h}_{ae,k} + \tilde{\mathbf{G}}_{ce,k}^H \phi^*)] \otimes (\mathbf{I}_M \phi)) \end{bmatrix} \\ &= \begin{bmatrix} \Sigma_{he,k}^{1/2} \mathbf{\Xi}_e (\mathbf{h}_{ae,k} + \tilde{\mathbf{G}}_{ce,k}^H \phi^*) \\ \Sigma_{ge,k}^{1/2T} (\mathbf{\Xi}_e \otimes \mathbf{I}_M) ((\mathbf{h}_{ae,k} + \tilde{\mathbf{G}}_{ce,k}^H \phi^*) \otimes \phi) \end{bmatrix}. \quad (42) \end{aligned}$$

The the step in ( $k$ ) in (42) is obtained by invoking the property  $(\mathbf{A} \otimes \mathbf{b})\mathbf{c} = (\mathbf{A}\mathbf{c}) \otimes \mathbf{b}$ , where  $\mathbf{b}$  and  $\mathbf{c}$  are column vectors. The expression of  $f_{1,k}^{\text{full}}$  can be rewritten as

$$\begin{aligned} f_{1,k}^{\text{full}} &= \Delta \mathbf{h}_{ae,k}^H \mathbf{\Xi}_e \Delta \mathbf{h}_{ae,k} + \text{Tr}\{(\mathbf{\Xi}_e \Delta \mathbf{G}_{ce,k}^H \mathbf{E} \Delta \mathbf{G}_{ce,k})\} \\ & \quad + 2\text{Re}\{\text{Tr}\{(\mathbf{\Xi}_e \Delta \mathbf{G}_{ce,k}^H \phi^* \Delta \mathbf{h}_{ae,k}^H)\}\} \\ &\stackrel{(l)}{=} \Delta \mathbf{h}_{ae,k}^H \mathbf{\Xi}_e \Delta \mathbf{h}_{ae,k} + \text{vec}^T(\Delta \mathbf{G}_{ce,k}) (\mathbf{\Xi}_e \otimes \mathbf{E}^T) \text{vec}(\Delta \mathbf{G}_{ce,k}^*) \\ & \quad + 2\text{Re}\{\text{vec}^T(\Delta \mathbf{h}_{ae,k}^H) (\mathbf{\Xi}_e \otimes \phi^H) \text{vec}(\Delta \mathbf{G}_{ce,k}^*)\} \\ &= \mathbf{v}_{he,k}^H \Sigma_{he,k}^{1/2} \mathbf{\Xi}_e \Sigma_{he,k}^{1/2} \mathbf{v}_{he,k} + \mathbf{v}_{ge,k}^T \Sigma_{ge,k}^{1/2T} (\mathbf{\Xi}_e \otimes \mathbf{E}^T) \Sigma_{ge,k}^{1/2*} \mathbf{v}_{ge,k}^* \\ & \quad + 2\text{Re}\{\mathbf{v}_{he,k}^H \Sigma_{he,k}^{1/2} (\mathbf{\Xi}_e \otimes \phi^H) \Sigma_{ge,k}^{1/2*} \mathbf{v}_{ge,k}^*\} \\ & \triangleq \tilde{\mathbf{v}}_{e,k}^H \tilde{\mathbf{A}}_{e,k} \tilde{\mathbf{v}}_{e,k}, \quad (43) \end{aligned}$$



where  $(l)$  is obtained by invoking  $\text{Tr}(\mathbf{ABCD}) = \text{vec}^T(\mathbf{D})(\mathbf{A} \otimes \mathbf{C}^T)\text{vec}(\mathbf{B}^T)$ , and

$$\tilde{\mathbf{A}}_{e,k} \triangleq \begin{bmatrix} \Sigma_{he,k}^{1/2} \Xi_e \Sigma_{he,k}^{1/2} & \Sigma_{he,k}^{1/2} (\Xi_e \otimes \phi^H) \Sigma_{ge,k}^{1/2*} \\ \Sigma_{ge,k}^{1/2T} (\Xi_e \otimes \phi) \Sigma_{he,k}^{1/2} & \Sigma_{ge,k}^{1/2T} (\Xi_e \otimes \mathbf{E}^T) \Sigma_{ge,k}^{1/2*} \end{bmatrix}. \quad (44)$$

Then, the outage probability constraint (40b) is recast as

$$\Pr_{\mathbf{g}_{he,k}, \mathbf{g}_{ge,k}} \left\{ \tilde{\mathbf{v}}_{e,k}^H \tilde{\mathbf{A}}_{e,k} \tilde{\mathbf{v}}_{e,k} + 2\text{Re}\{\tilde{\mathbf{u}}_{e,k}^H \tilde{\mathbf{v}}_{e,k}\} + c_k^{\text{full}} \geq 0 \right\} \geq 1 - \rho_k, \forall k \in \mathcal{K}. \quad (45)$$

It is seen from (39)-(45) that Lemma 1 can be applied as long as the CSI error  $\mathbf{g}_{ge,k}$  and  $\mathbf{g}_{he,k}$  follow the CSCG distribution and regardless of the correlation property of CSI errors. By leveraging Lemma 1 again, the chance constraints in (39) are approximated by computable constraints, which results the equivalent OCR-PM problem as

$$\min_{\mathbf{W}, \mathbf{Z}, \tilde{\phi}, \tilde{\mathbf{x}}, \tilde{\mathbf{y}}} \text{Tr}(\mathbf{W} + \mathbf{Z}) \quad (46a)$$

$$\text{s.t. } \text{Tr}(\tilde{\mathbf{A}}_{e,k}) - \sqrt{-2 \ln(\rho_k)} \cdot \tilde{x}_k + \ln(\rho_k) \cdot \tilde{y}_k + c_k^{\text{full}} \geq 0, \quad (46b)$$

$$\left\| \begin{bmatrix} \text{vec}(\tilde{\mathbf{A}}_{e,k}) \\ \sqrt{2} \tilde{\mathbf{u}}_{e,k} \end{bmatrix} \right\|_2 \leq \tilde{x}_k, \quad (46c)$$

$$\tilde{y}_k \mathbf{I}_{MN_t + N_t} + \tilde{\mathbf{A}}_{e,k} \geq \mathbf{0}, \tilde{y}_k \geq 0, \quad (46d)$$

$$(10b), (9d), (9e), (9f), (9g), \quad (46e)$$

where  $\forall k \in \mathcal{K}$ ,  $\tilde{\mathbf{x}} = [\tilde{x}_1, \tilde{x}_2, \dots, \tilde{x}_K]^T$  and  $\tilde{\mathbf{y}} = [\tilde{y}_1, \tilde{y}_2, \dots, \tilde{y}_K]^T$  are introduced variables.

2) *Optimization of Transmit Beamforming and AN*: The AO method is also utilized to decouple the variables. When  $\phi$  is fixed, the  $\tilde{\mathbf{A}}_{e,k}$  and  $\tilde{\mathbf{u}}_{e,k}$  are linear functions of  $\mathbf{W}$  and  $\mathbf{Z}$ . Then, by removing  $\text{rank}(\mathbf{W}) = 1$ , the  $\mathbf{W}$  and  $\mathbf{Z}$  can be obtained by the SDR, and the corresponding optimization problem is

$$\min_{\mathbf{W}, \mathbf{Z}, \tilde{\mathbf{x}}, \tilde{\mathbf{y}}} \text{Tr}(\mathbf{W} + \mathbf{Z}) \quad (47a)$$

$$\text{s.t. } (46b), (46c), (46d), (10b), (9d), (9e), \forall k \in \mathcal{K}. \quad (47b)$$

The SDR in Problem (47) is tight. Theorem 1 still holds, and the proofs are the same as those under partial CSI errors.

3) *Optimization of Phase Shifts at the IRS*: When  $\mathbf{W}$  and  $\mathbf{Z}$  are fixed, the optimization problem for  $\phi$  becomes a feasibility check problem, and slack variables  $\delta = [\delta_0, \delta_1, \delta_2, \dots, \delta_K]^T$  are introduced to improve the AO's convergence. Due to the existence of the direct communication link in this paper, the optimization variable can only be  $\phi$  instead of  $\mathbf{E}$ , and almost all constraints are nonconvex w.r.t.  $\phi$ . Thus, we propose novel mathematical manipulations to transform Problem (46) into a convex one w.r.t.  $\phi$  in the following three steps. It is noted that these manipulations are different but inherited from those in the scenario of partial CSI errors.

**Step 1**: Transform the constraint in (46b) to be convex.

We find that  $c_k^{\text{full}} = c_k^{\text{partial}}$ , thus the manipulation on  $c_k^{\text{full}}$  is the same as that on  $c_k^{\text{partial}}$  under partial CSI errors. Then we have  $c_k^{\text{full}} = (23)$ , the expression of which is concave w.r.t.  $\phi$ . The  $\text{Tr}[\tilde{\mathbf{A}}_{e,k}]$  are transformed as

$$\begin{aligned} \text{Tr}[\tilde{\mathbf{A}}_{e,k}] &= \text{Tr}[\Sigma_{he,k}^{1/2} \Xi_e \Sigma_{he,k}^{1/2}] + \left\{ \text{Tr}[\Sigma_{ge,k}^{1/2T} (\Xi_e \otimes \mathbf{E}^T) \Sigma_{ge,k}^{1/2*}] \right\}^* \\ &\stackrel{(m)}{=} \text{Tr}[\Sigma_{he,k}^{1/2} \Xi_e \Sigma_{he,k}^{1/2}] + \text{Tr}[\Sigma_{ge,k}^{1/2H} (\Xi_e^T \otimes \mathbf{E}) \Sigma_{ge,k}^{1/2}] \\ &\stackrel{(n)}{=} \text{Tr}[\Sigma_{he,k}^{1/2} \Xi_e \Sigma_{he,k}^{1/2}] + \phi^T \Sigma_{g_{x,k}} (\mathbf{I}_{MN_t} \otimes \Xi_e) \Sigma_{g_{x,k}}^H \phi^*, \end{aligned} \quad (48)$$

where the second term of  $(m)$  is the same as the  $\text{Tr}[\mathbf{A}_{e,k}]$  under partial CSI errors, and the second term of  $(n)$  is obtained in the same way of (24). The second term in  $(n)$  of (48) is equivalently transformed to be concave w.r.t.  $\phi$  as in (25). By substituting (25) into (48), we reformulate (48) into a concave form as

$$\begin{aligned} \text{Tr}[\tilde{\mathbf{A}}_{e,k}] &= \text{Tr}[\Sigma_{he,k}^{1/2} \Xi_e \Sigma_{he,k}^{1/2}] + \phi^T [\Sigma_{g_{x,k}} (\mathbf{I}_{MN_t} \otimes \Xi_e) \Sigma_{g_{x,k}}^H \\ &\quad - c_{t,k} \mathbf{I}_M] \phi^* + M c_{t,k}. \end{aligned} \quad (49)$$

By substituting (23) and (49) into the constraint in (46b), we can obtain a convex constraint equivalent to (46b) as

$$(46b) \Leftrightarrow c_{tc,k}^{\text{full}}(\phi) \geq 0, \quad (50)$$

which is convex, and

$$\begin{aligned} c_{tc,k}^{\text{full}}(\phi) &\triangleq \text{Tr}[\Sigma_{he,k}^{1/2} \Xi_e \Sigma_{he,k}^{1/2}] + \phi^T [\Sigma_{g_{x,k}} (\mathbf{I}_{MN_t} \otimes \Xi_e) \Sigma_{g_{x,k}}^H \\ &\quad - c_{t,k} \mathbf{I}_M] \phi^* + M c_{t,k} - \sqrt{-2 \ln(\rho_k)} \cdot x_k + \ln(\rho_k) \cdot y_k \\ &\quad + \phi^T (\tilde{\mathbf{G}}_{ce,k} \Xi_e \tilde{\mathbf{G}}_{ce,k}^H - c_{c,k} \mathbf{I}_M) \phi^* + 2\text{Re}\{\phi^T \tilde{\mathbf{G}}_{ce,k} \Xi_e \mathbf{h}_{ae,k}\} \\ &\quad + \mathbf{h}_{ae,k}^H \Xi_e \mathbf{h}_{ae,k} + M c_{c,k} + \tilde{\sigma}_{e,k}^2 - \delta_k. \end{aligned} \quad (51)$$

**Step 2**: Transform the constraint in (46c) to be convex.

We can find an upper bound of the left hand side of constraint (46c) as follows:

$$\begin{aligned} \|\text{vec}(\tilde{\mathbf{A}}_{e,k})\|^2 &= \left\| \Sigma_{he,k}^{1/2} \Xi_e \Sigma_{he,k}^{1/2} \right\|^2 + 2 \left\| \Sigma_{ge,k}^{1/2T} (\Xi_e \otimes \phi) \Sigma_{he,k}^{1/2} \right\|^2 \\ &\quad + \left\| \Sigma_{ge,k}^{1/2T} (\Xi_e \otimes \mathbf{E}^T) \Sigma_{ge,k}^{1/2*} \right\|^2 \end{aligned} \quad (52a)$$

$$\stackrel{(o)}{=} \left\| (\Sigma_{he,k}^{1/2T} \otimes \Sigma_{he,k}^{1/2}) \text{vec}(\Xi_e) \right\|^2 + 2 \left\| (\Sigma_{he,k}^{1/2T} \otimes \Sigma_{ge,k}^{1/2T}) \text{vec}(\Xi_e \otimes \phi) \right\|^2 \\ + \left\| (\Sigma_{ge,k}^{1/2} \otimes \Sigma_{ge,k}^{1/2T}) \text{vec}(\Xi_e \otimes \mathbf{E}^T) \right\|^2 \quad (52b)$$

$$\stackrel{(p)}{\leq} \sigma_{\max}^2(\Sigma_{he,k}^{1/2T} \otimes \Sigma_{he,k}^{1/2}) \|\Xi_e\|_F^2 + 2\sigma_{\max}^2(\Sigma_{he,k}^{1/2T} \otimes \Sigma_{ge,k}^{1/2T}) M \|\Xi_e\|_F^2 \\ + \sigma_{\max}^2(\Sigma_{ge,k}^{1/2} \otimes \Sigma_{ge,k}^{1/2T}) M^2 \|\Xi_e\|_F^2 \quad (52c)$$

$$\stackrel{(q)}{=} [\lambda_{\max}(\Sigma_{he,k}) + \lambda_{\max}(\Sigma_{ge,k}) M]^2 \|\Xi_e\|_F^2, \quad (52d)$$

$$\|\tilde{\mathbf{u}}_{e,k}\|^2 = \left\| \begin{bmatrix} \Sigma_{he,k}^{1/2} \Xi_e (\mathbf{h}_{ae,k} + \tilde{\mathbf{G}}_{ce,k}^H \phi^*) \\ \Sigma_{ge,k}^{1/2T} (\Xi_e \otimes \mathbf{I}_M) ((\mathbf{h}_{ae,k} + \tilde{\mathbf{G}}_{ce,k}^H \phi^*) \otimes \phi) \end{bmatrix} \right\|^2 \quad (53a)$$

$$= \left\| \Sigma_{he,k}^{1/2} \Xi_e (\mathbf{h}_{ae,k} + \tilde{\mathbf{G}}_{ce,k}^H \phi^*) \right\|^2 \\ + \left\| \Sigma_{ge,k}^{1/2T} (\Xi_e \otimes \mathbf{I}_M) ((\mathbf{h}_{ae,k} + \tilde{\mathbf{G}}_{ce,k}^H \phi^*) \otimes \phi) \right\|^2 \quad (53b)$$

$$\stackrel{(r)}{\leq} \sigma_{\max}^2(\Sigma_{he,k}^{1/2} \Xi_e) \|\mathbf{h}_{ae,k} + \tilde{\mathbf{G}}_{ce,k}^H \phi^*\|^2 \\ + \sigma_{\max}^2(\Sigma_{ge,k}^{1/2T} (\Xi_e \otimes \mathbf{I}_M)) M \|\mathbf{h}_{ae,k} + \tilde{\mathbf{G}}_{ce,k}^H \phi^*\|^2 \quad (53c)$$

$$= [\lambda_{\max}^2(\Sigma_{he,k}^{1/2} \Xi_e) + \lambda_{\max}^2(\Sigma_{ge,k}^{1/2T} (\Xi_e \otimes \mathbf{I}_M)) M] \\ \cdot \|\mathbf{h}_{ae,k} + \tilde{\mathbf{G}}_{ce,k}^H \phi^*\|^2, \quad (53d)$$

where the  $(o)$  is obtained by using Lemma 3 as follows.

**Lemma 3.** For any three matrices  $\mathbf{O}$ ,  $\mathbf{P}$ , and  $\mathbf{Q}$ , the equality  $\|\mathbf{POQ}\|^2 = \|(\mathbf{Q}^T \otimes \mathbf{P}) \text{vec}(\mathbf{O})\|^2$  holds.

**Proof.** Please refer to Appendix C.  $\square$

The  $(p)$  and  $(r)$  are obtained by using  $\|\mathbf{Ax}\|_2 \leq \|\mathbf{A}\|_2 \|\mathbf{x}\|_2$ ,  $\|\mathbf{A}\|_2 = \sigma_{\max}\{\mathbf{A}\}$ . The third term of (52c) is obtained due



to the property in (29). When  $\Sigma_{he,k}^{1/2T} \otimes \Sigma_{he,k}^{1/2}$ ,  $\Sigma_{he,k}^{1/2T} \otimes \Sigma_{ge,k}^{1/2T}$  and  $\Sigma_{ge,k}^{1/2} \otimes \Sigma_{ge,k}^{1/2T}$  are unitary matrices, the inequality of (p) in (52c) becomes an equality. The (q) is obtained due to  $\text{eig}(\mathbf{A} \otimes \mathbf{B}) = \text{eig}(\mathbf{A}) \otimes \text{eig}(\mathbf{B})$  and  $\sigma_{\max}\{\mathbf{A}\} = \lambda_{\max}\{\mathbf{A}\}$  when  $\mathbf{A} \in \mathbb{H}^n$ . The second term of (53c) is obtained due to Lemma 2. When  $\Sigma_{he,k}^{1/2} \Xi_e$  and  $\Sigma_{ge,k}^{1/2T}(\Xi_e \otimes \mathbf{I}_M)$  are unitary matrices, the inequality of (r) in (53c) becomes an equality. By using the upper bounds in (52d) and (53d), the original constraint (46c) can be approximated as

$$\left\| \frac{[\lambda_{\max}(\Sigma_{he,k}) + \lambda_{\max}(\Sigma_{ge,k})M] \text{vec}(\Xi_e)}{\sqrt{2[\lambda_{\max}^2(\Sigma_{he,k}^{1/2} \Xi_e) + \lambda_{\max}^2(\Sigma_{ge,k}^{1/2T}(\Xi_e \otimes \mathbf{I}_M))]M}} (\mathbf{h}_{ae,k} + \tilde{\mathbf{G}}_{ce,k}^H \phi) \right\| \leq \tilde{x}_k. \quad (54)$$

It is readily to check that when (54) holds, the (46c) always holds, i.e., (46c)  $\Leftrightarrow$  (54). When the inequalities of (p) and (r) both become equalities, we have (46c)  $\Leftrightarrow$  (54).

**Step3:** Transform the constraint in (46d) to be convex.

The transformation is similar as the process in (32d). By moving  $\tilde{\mathbf{A}}_{e,k}$  in (46d) from the left hand side to the right hand side of the inequality, we have

$$\tilde{y}_k \mathbf{I}_{MN_t+N_t} \geq -\tilde{\mathbf{A}}_{e,k}. \quad (55)$$

Then, we define a matrix  $\mathbf{B}_{e,k}$  as

$$\mathbf{B}_{e,k} \triangleq \begin{bmatrix} \Sigma_{he,k}^{1/2} (c_g \mathbf{I}_{N_t}) \Sigma_{he,k}^{1/2} & \Sigma_{he,k}^{1/2} (c_g \mathbf{I}_{N_t} \otimes \phi^H) \Sigma_{ge,k}^{1/2*} \\ \Sigma_{ge,k}^{1/2T} (c_g \mathbf{I}_{N_t} \otimes \phi) \Sigma_{he,k}^{1/2} & \Sigma_{ge,k}^{1/2T} (c_g \mathbf{I}_{N_t} \otimes \mathbf{E}^T) \Sigma_{ge,k}^{1/2*} \end{bmatrix}. \quad (56)$$

By adding the matrix  $\mathbf{B}_{e,k}$  on both sides of the inequality in (55), we have the following equivalent inequality.

$$\begin{aligned} & y_k \mathbf{I}_{(MN_t+N_t)} + \mathbf{B}_{e,k} \geq \\ & \begin{bmatrix} \Sigma_{he,k}^{1/2} (c_g \mathbf{I}_{N_t}) \Sigma_{he,k}^{1/2} & \Sigma_{he,k}^{1/2} (c_g \mathbf{I}_{N_t} \otimes \phi^H) \Sigma_{ge,k}^{1/2*} \\ \Sigma_{ge,k}^{1/2T} (c_g \mathbf{I}_{N_t} \otimes \phi) \Sigma_{he,k}^{1/2} & \Sigma_{ge,k}^{1/2T} (c_g \mathbf{I}_{N_t} \otimes \mathbf{E}^T) \Sigma_{ge,k}^{1/2*} \end{bmatrix} - \tilde{\mathbf{A}}_{e,k} \\ \Leftrightarrow & y_k \mathbf{I}_{(MN_t+N_t)} + \mathbf{B}_{e,k} \geq \\ & \begin{bmatrix} \Sigma_{he,k}^{1/2} (\mathbf{V}_{wgz} \mathbf{V}_{wgz}^H) \Sigma_{he,k}^{1/2} & \Sigma_{he,k}^{1/2} ((\mathbf{V}_{wgz} \mathbf{V}_{wgz}^H) \otimes \phi^H) \Sigma_{ge,k}^{1/2*} \\ \Sigma_{ge,k}^{1/2T} ((\mathbf{V}_{wgz} \mathbf{V}_{wgz}^H) \otimes \phi) \Sigma_{he,k}^{1/2} & \Sigma_{ge,k}^{1/2T} ((\mathbf{V}_{wgz} \mathbf{V}_{wgz}^H) \otimes \mathbf{E}^T) \Sigma_{ge,k}^{1/2*} \end{bmatrix} \\ \Leftrightarrow & y_k \mathbf{I}_{(MN_t+N_t)} + \mathbf{B}_{e,k} \geq \\ & \begin{bmatrix} \Sigma_{he,k}^{1/2} \mathbf{V}_{wgz} \\ \Sigma_{ge,k}^{1/2T} (\mathbf{V}_{wgz} \otimes \phi) \end{bmatrix} \begin{bmatrix} \mathbf{V}_{wgz}^H \Sigma_{he,k}^{1/2} & (\mathbf{V}_{wgz}^H \otimes \phi^H) \Sigma_{ge,k}^{1/2*} \end{bmatrix}. \quad (57) \end{aligned}$$

Similarly, we find that  $\lambda_{\min}\{\mathbf{B}_{e,k}\} = 0$ , which is proved in Appendix D. Thus, from (57), we have the convex linear matrix inequality (LMI) constraint as

$$(46d) \Leftrightarrow \sqrt{\tilde{y}_k} \stackrel{(s)}{\geq} \left\| \begin{bmatrix} \Sigma_{he,k}^{1/2} \mathbf{V}_{wgz} \\ \Sigma_{ge,k}^{1/2T} (\mathbf{V}_{wgz} \otimes \phi) \end{bmatrix} \right\|, \tilde{y}_k \geq 0, \quad (58)$$

where (s) is obtained due to  $\|\mathbf{A}\|_2 = \sqrt{\max \text{eig}(\mathbf{A}^H \mathbf{A})}$ .

Finally, based on the mathematical manipulations of **Steps 1-3** above and by using the CCP method [27], the problem for optimizing  $\phi$  can be reformulated as

$$\max_{\phi, \delta, \tilde{x}, \tilde{y}} \sum_{k=0}^K \delta_k - \lambda^{(t+1)} \sum_{l=1}^{2M} b_l \quad (59a)$$

$$\text{s.t. (50), (54), (58), (35), (22e), (38c), (38d), (38e).} \quad (59b)$$

The same techniques under partial CSI errors can be utilized here to solve Problem (59), which is omitted to make the paper compact.

## C. Overall Algorithm, Convergence and Complexity Analysis

1) *Overall Algorithm:* The overall AO algorithm proposed under partial/full CSI errors is summarized in Algorithm 1, where the expression of ‘‘Problem (A)/Problem (B)’’ is utilized to denote that the Problem (A) under partial CSI errors can be replaced by the Problem (B) under full CSI errors. By iteratively solving Problem (20)/ Problem (47) and Problem (38)/ Problem (59) optimally in Step 3 and Step 4 in Algorithm 1, the transmit power can be monotonically reduced with guaranteed convergence. We prove the convergence of the proposed AO algorithm as follows.

### Algorithm 1 Alternating Optimization Algorithm

- 1) **Parameter Setting:** Set the maximum number of iterations  $n_{\max}$  and the first iterative number  $n = 1$ ; Give the error tolerance  $\varepsilon$ ;
- 2) **Initialize** the variables  $\mathbf{w}^{(1)}$ ,  $\mathbf{Z}^{(1)}$  and  $\phi^{(1)}$  in the feasible region; Compute the OF value of Problem (19)/ Problem (46) (i.e., the transmit power) as  $p(\mathbf{w}^{(1)}, \mathbf{Z}^{(1)})$ ;
- 3) **Solve** Problem (20)/ Problem (47) to obtain the  $\mathbf{w}^{(n)}$ ,  $\mathbf{Z}^{(n)}$  by fixing  $\phi^{(n)}$ ; Calculate the OF value of Problem (20)/ Problem (47) as  $p(\mathbf{w}^{(n)}, \mathbf{Z}^{(n)})$ ;
- 4) **Solve** Problem (38)/ Problem (59) to obtain  $\phi^{(n+1)}$  by fixing  $\mathbf{w}^{(n)}$ ,  $\mathbf{Z}^{(n)}$ ;
- 5) **If**  $|p(\mathbf{w}^{(n+1)}, \mathbf{Z}^{(n+1)}) - p(\mathbf{w}^{(n)}, \mathbf{Z}^{(n)})| / p(\mathbf{w}^{(n)}, \mathbf{Z}^{(n)}) < \varepsilon$  or  $n \geq n_{\max}$ , terminate. Otherwise, update  $n \leftarrow n+1$  and jump to Step 3.

2) *Convergence:* Denote the OF value of Problem (19)/ Problem (46) and Problem (20)/ Problem (47) (i.e., the transmit power) with a feasible solution  $(\phi, \mathbf{w}, \mathbf{Z})$  as  $p(\phi, \mathbf{w}, \mathbf{Z})$ . As shown in step 4 of Algorithm 1, if there exists a feasible solution to Problem (38)/ Problem (59), i.e.,  $(\phi^{(n+1)}, \mathbf{w}^{(n)}, \mathbf{Z}^{(n)})$  exists, it is also feasible to Problem (20)/ Problem (47). Then,  $(\phi^{(n+1)}, \mathbf{w}^{(n+1)}, \mathbf{Z}^{(n+1)})$  and  $(\phi^{(n)}, \mathbf{w}^{(n)}, \mathbf{Z}^{(n)})$  in step 3 are the feasible solutions to Problem (20)/ Problem (47) in the  $(n+1)$ th and  $(n)$ th iterations, respectively. It then follows that  $p(\phi^{(n+1)}, \mathbf{w}^{(n+1)}, \mathbf{Z}^{(n+1)}) \stackrel{(u)}{\leq} p(\phi^{(n+1)}, \mathbf{w}^{(n)}, \mathbf{Z}^{(n)}) \stackrel{(v)}{=} p(\phi^{(n)}, \mathbf{w}^{(n)}, \mathbf{Z}^{(n)})$ , where (u) holds because for given  $\phi^{(n+1)}$  in step 3 of Algorithm 1,  $\mathbf{w}^{(n+1)}, \mathbf{Z}^{(n+1)}$  is the optimal solution to Problem (20)/ Problem (47); and (v) holds because the OF of Problem (20)/ Problem (47) is regardless of  $\phi$  and only depends on  $\mathbf{w}, \mathbf{Z}$ .

3) *Complexity:* The computational complexity of all the resulted convex problems in Algorithm 1 can be measured in terms of their worst-case runtime by counting the complexity of LMI and second-order cone (SOC) constraints and ignoring the complexity of linear constraints, and the general expression for complexity has been given in [27]. For Problem (20), the number of variables is  $n_1 = 2N_t^2$ . The number of LMIs in (19d) is  $K$  with the size of  $MN_t$ . The number of LMIs in (9d) and (9e) is 2 with the size of  $N_t$ . The number of SOC in (19c) is  $K$  with the size of  $M^2 N_t^2 + MN_t$ . Thus, the approximate complexity of Problem (20) is

$$C_{\text{WZ}} = \mathcal{O}\{[KMN_t + 2N_t + 2K]^{1/2} n_1 [n_1^2 + n_1(KM^2 + 2)N_t^2 + (KM^3 + 2)N_t^3 + n_1 K(M^2 N_t^2 + MN_t)^2]\}. \quad (60)$$

For Problem (38), the number of variables is  $n_2 = M$ . The number of LMIs in (31) is  $K$  with the size of  $N_t^2 + N_t$ . The number of SOC in (33) is  $K$  with the size of  $MN_t^2$ . Thus, the approximate complexity of Problem (38) is

$$C_\phi = O\{(4K)^{1/2}n_2[n_2^2 + n_2(K(N_t^2 + N_t)^2 + KM^2N_t^4)]\}. \quad (61)$$

Altogether, the approximate computational complexity in each iteration under partial CSI errors is  $C_{\mathbf{WZ}} + C_\phi$ .

Similarly, for Problem (47), the number of variables is  $\bar{n}_1 = 2N_t^2$ . The number of SOC in (46c) is  $K$  with the size of  $(MN_t + N_t)^2 + MN_t + N_t$ . The number of LMIs in (46d) is  $K$  with the size of  $MN_t + N_t$ . The number of LMIs in (9d) and (9e) is 2 with the size of  $N_t$ . Thus, the approximate complexity of Problem (47) is

$$\begin{aligned} \tilde{C}_{\mathbf{WZ}} = & O\{[K(MN_t + N_t) + 2N_t + 2K]^{1/2}\bar{n}_1[\bar{n}_1^2 \\ & + \bar{n}_1(K(MN_t + N_t)^2 + 2N_t^2) + (K(MN_t + N_t)^3 + 2N_t^3) \\ & + \bar{n}_1K((MN_t + N_t)^2 + MN_t + N_t)^2]\}. \end{aligned} \quad (62)$$

For Problem (59), the number of variables is  $\bar{n}_2 = M$ . The number of SOC in (54) is  $K$  with the size of  $N_t^2 + N_t$ . The number of LMIs in (58) is  $K$  with the size of  $N_t^2 + N_t$ . Thus, the approximate complexity of Problem (59) is

$$\begin{aligned} \tilde{C}_\phi = & O\{(4K)^{1/2}\bar{n}_2[\bar{n}_2^2 + \bar{n}_2(K(N_t^2 + N_t)^2 \\ & + K(MN_t^2 + N_t^2)^2)]\}. \end{aligned} \quad (63)$$

Altogether, the approximate computational complexity in each iteration under full CSI errors is  $\tilde{C}_{\mathbf{WZ}} + \tilde{C}_\phi$ .

#### D. Extensions to the Multiple-Bob Case

Consider a multiple-Bob extension where  $K$  legitimate users (Bobs) are served by Alice, and are eavesdropped by  $K$  Eves. Specifically, the  $k$ th Bob is eavesdropped by the  $k$ th Eve [36]. Then, the transmit signal is revised into

$$\mathbf{x} = \bar{\mathbf{W}}\bar{\mathbf{s}} + \mathbf{z} = \sum_{k=1}^K \mathbf{w}_k s_k + \mathbf{z}, \quad (64)$$

where  $\bar{\mathbf{s}} = [s_1, s_2, \dots, s_K]^T$ , and  $s_k$  is the data symbol intended for the  $k$ th Bob.  $\bar{\mathbf{W}} = [\mathbf{w}_1, \mathbf{w}_2, \dots, \mathbf{w}_K]$  contains all transmit beamforming vectors, and  $\mathbf{w}_k$  is the beamforming vector for the  $k$ th Bob. We denote the channel vectors spanning from the BS to the  $k$ th Bob and from the RIS to the  $k$ th Bob by  $\mathbf{h}_{ab,k} \in N_t \times 1$  and  $\mathbf{h}_{rb,k} \in N_t \times 1$ , respectively. The received signal at the  $k$ th Bob are expressed as

$$y_{b,k} = \hat{\mathbf{h}}_{b,k}^H \mathbf{x} + n_{b,k} = (\mathbf{h}_{ab,k}^H + \mathbf{h}_{rb,k}^H \Phi \mathbf{G}_{ar}) \mathbf{x} + n_{b,k}, \quad (65)$$

where  $\forall k \in \mathcal{K}$ ,  $\hat{\mathbf{h}}_{b,k}^H \triangleq \mathbf{h}_{ab,k}^H + \phi^T \mathbf{G}_{cb,k}$  is the overall channel from Alice to Bob, and  $\mathbf{G}_{cb,k} \triangleq \text{diag}(\mathbf{h}_{rb,k}^H) \mathbf{G}_{ar}$  is the cascaded channel from Alice to Bob via IRS. The  $n_{b,k} \sim \mathcal{CN}(0, \sigma_{b,k}^2)$  is the AWGN received at the  $k$ th Bob.

For security provisioning, we make a worst-case assumption regarding the capabilities of the potential eavesdroppers, and assume that Eves are able to cancel all multiuser interference before decoding the information transmitted to a given Bob [36]. Then the achievable data rate of the  $k$ th Eve can be written as

$$C_{e,k}(\mathbf{W}_k, \mathbf{Z}, \Phi) = \log[1 + |\hat{\mathbf{h}}_{e,k}^H \mathbf{w}_k|^2 / (\sigma_{e,k}^2 + |\hat{\mathbf{h}}_{e,k}^H \mathbf{z}|^2)], \quad (66)$$

where  $\mathbf{W}_k = \mathbf{w}_k \mathbf{w}_k^H$ . To ensure a fair condition for Bob and Eve, we assume that Bobs can also cancel the multiuser interference, thus the achievable data rate of the  $k$ th Bob is

$$C_{b,k}(\mathbf{W}_k, \mathbf{Z}, \Phi) = \log[1 + |\hat{\mathbf{h}}_{b,k}^H \mathbf{w}_k|^2 / (\sigma_{b,k}^2 + |\hat{\mathbf{h}}_{b,k}^H \mathbf{z}|^2)]. \quad (67)$$

1) *Scenario 1: Partial CSI Errors:* By using (66) and (67), the OCR-PM problem with partial CSI errors is formulated as

$$\min_{\{\mathbf{W}_k\}_{k=1}^K, \mathbf{Z}, \Phi} \sum_{k=1}^K \text{Tr}(\mathbf{W}_k) + \text{Tr}(\mathbf{Z}) \quad (68a)$$

$$\text{s.t. } C_{b,k}(\mathbf{W}_k, \mathbf{Z}, \Phi) \geq \log \gamma_k, \quad (68b)$$

$$\text{Pr}_{\mathbf{g}_{e,k}} \{C_{e,k}(\mathbf{W}_k, \mathbf{Z}, \Phi) \leq \log \beta_k\} \geq 1 - \rho_k, \quad (68c)$$

$$\mathbf{Z} \geq 0, \quad (68d)$$

$$\mathbf{W}_k \geq 0, \quad (68e)$$

$$\text{rank}(\mathbf{W}_k) = 1, \quad (68f)$$

$$|\phi_m| = 1, m = 1, \dots, M, \quad (68g)$$

where  $\forall k \in \mathcal{K}$ ,  $\gamma_k \geq 1$  and  $\beta_k \geq 1$  are constant values, and  $\gamma_k \geq \beta_k$  is imposed to ensure a positive secrecy rate. Obviously, the constraints of (68b)-(68g) above are similar to the constraints of (9b)-(9g) above, thus the problem reformulation and problem solving can be extended directly from those for single-Bob case in Section III-A above. The main extensions or modifications from the single-Bob case lie in the following aspects. (1) The beamforming vector to be optimized is extended from a single  $\mathbf{w}$  to multiple  $\{\mathbf{w}_k\}_{k=1}^K$ . (2) When optimizing the transmit beamforming vectors and AN similar as in Problem (20), the  $\Xi_e$  in constraints (19b),(19c),(19d) should be modified into  $\Xi_{e,k} \triangleq (\beta_k - 1)\mathbf{Z} - \mathbf{W}_k$ , and the constraint (10b) should be modified into

$$\hat{\mathbf{h}}_{b,k}^H [\mathbf{W}_k - (\gamma_k - 1)\mathbf{Z}] \hat{\mathbf{h}}_{b,k} \geq (\gamma_k - 1)\sigma_{b,k}^2, \forall k \in \mathcal{K}. \quad (69)$$

(3) When optimizing the phase shifts at the IRS similar as in Problem (38), the  $\Xi_e$  in constraints (26),(31),(33) should be modified into  $\Xi_{e,k}$ , and the constraint (35) should be modified into

$$\begin{aligned} & (\mathbf{h}_{ab,k}^H + \phi^T \mathbf{G}_{cb,k}) [\mathbf{W}_k - (\gamma_k - 1)\mathbf{Z}] (\mathbf{h}_{ab,k} + \mathbf{G}_{cb,k}^H \phi^*) \\ & \geq \tilde{\sigma}_{b,k}^2 + \delta_{K+k}, \forall k \in \mathcal{K} \Leftrightarrow \\ & \phi^T [\mathbf{G}_{cb,k} \Xi_{b,k} \mathbf{G}_{cb,k}^H - c_{b,k} \mathbf{I}_M] \phi^* + 2\text{Re}\{\phi^T \mathbf{G}_{cb,k} \Xi_{b,k} \mathbf{h}_{ab,k}\} \\ & + \mathbf{h}_{ab,k}^H \Xi_{b,k} \mathbf{h}_{ab,k} - \tilde{\sigma}_{b,k}^2 - \delta_{K+k} + M c_{b,k} \geq 0, \forall k \in \mathcal{K}, \end{aligned} \quad (70)$$

where  $\tilde{\sigma}_{b,k}^2 \triangleq (\gamma_k - 1)\sigma_{b,k}^2$ ,  $\Xi_{b,k} = \mathbf{W}_k - (\gamma_k - 1)\mathbf{Z}$ ,  $c_{b,k} = \lambda_{\max}\{\mathbf{G}_{cb,k} \mathbf{W}_k \mathbf{G}_{cb,k}^H\}$ , and  $\delta_{K+k}$  is the slack variable for the  $k$ th Bob. The slack variables in  $\delta = [\delta_0, \delta_1, \delta_2, \dots, \delta_K]^T$  are increased into  $\delta = [\delta_1, \delta_2, \dots, \delta_{2K}]^T$ . Due to the page limit, the detailed derivations for solving Problem (68) are omitted, which can be found in our full journal version in [37].

2) *Scenario 2: Full CSI Errors:* In the scenario of full CSI errors, the OCR-PM problem is formulated as

$$\min_{\{\mathbf{W}_k\}_{k=1}^K, \mathbf{Z}, \Phi} \sum_{k=1}^K \text{Tr}(\mathbf{W}_k) + \text{Tr}(\mathbf{Z}) \quad (71a)$$

$$\text{Pr}_{\mathbf{g}_{e,k}, \mathbf{g}_{e,k}} \{C_{e,k}(\mathbf{W}_k, \mathbf{Z}, \Phi) \leq \log \beta_k\} \geq 1 - \rho_k, \quad (71b)$$

$$(68b), (68d), (68e), (68f), (68g), \quad (71c)$$

where  $\forall k \in \mathcal{K}$ . Similarly, the reformulation and solving of Problem (71) can be extended directly from those for the

single-Bob case in Section III-B. The main extensions or modifications from the single-Bob case in the scenario of full CSI errors are the same as those in the scenario of partial CSI errors illustrated above. The detailed derivations for solving Problem (71) can be found in our full journal version in [37].

#### IV. SIMULATION RESULTS

Fig. 2 describes the considered IRS-aided secure communication system, where the Alice, IRS, and Bob are located [16] at (5,0,20) m, (0,50,2) m, and (3,50,0) m respectively. The Eves are randomly distributed on the line from (2,45,0) m to (2,55,0) m. We assume that the channel from Alice to the IRS is Rician fading and can be modeled as

$$\mathbf{G}_{ar} = \sqrt{L_0 d_{ar}^{-\alpha_{ar}}} \left( \sqrt{\frac{\beta_{ar}}{1+\beta_{ar}}} \mathbf{G}_{ar}^{LOS} + \sqrt{\frac{1}{1+\beta_{ar}}} \mathbf{G}_{ar}^{NLOS} \right), \quad (72)$$

where the pathloss at the reference distance is set to be  $L_0 = -40$  dB based on the 3GPP UMi model [38]. The distance from Alice to the IRS is denoted by  $d_{ar}$ , and the path loss exponent of the Alice-IRS link is denoted by  $\alpha_{ar}$ .  $\beta_{ar}$  is the corresponding Rician factor. The channel  $\mathbf{G}_{ar}$  contains the line of sight (LoS)  $\mathbf{G}_{ar}^{LOS}$  and the Rayleigh fading non-LoS (NLoS)  $\mathbf{G}_{ar}^{NLOS}$  components. Other channels are modeled similarly.

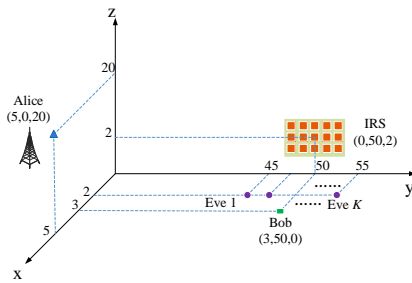


Fig. 2. Simulation setup of the considered IRS-aided secure communication

The path loss exponent for Alice-IRS channel is  $\alpha_{ar} = 2.2$  [27]. The path loss exponents for IRS-Bob channel  $\alpha_{rb}$  and IRS-Eve channel  $\alpha_{re}$  are 2 [27]. A more scattering environment is assumed for the direct links, thus the path loss exponents for Alice-Bob channel and Alice-Eve channel are  $\alpha_{ab} = \alpha_{ae} = 3.5$  [39]. The Rician factors for the Alice-IRS channel  $\beta_{ar}$ , IRS-Bob channel  $\beta_{rb}$ , IRS-Eve channel  $\beta_{re}$ , Alice-Bob channel  $\beta_{ab}$ , and Alice-Eve channel  $\beta_{ae}$  are 5 [26]. The Eves' outage probability is  $\rho_k = 0.05$  [27]. The noise power at Bob and Eves is set as  $-85$  dBm. The threshold for convergence is  $\varepsilon = 10^{-3}$  [26].

For comparison, we exploit the following baseline schemes.

- MRT-isoAN-randIRS: The maximum ratio transmission (MRT) based beamforming is performed, where  $\mathbf{w} = \sqrt{p_w} \frac{\hat{\mathbf{h}}_b}{\|\hat{\mathbf{h}}_b\|}$ , and  $p_w$  is the power allocated to Bob. The isotropic AN [35] is generated, where the AN covariance matrix is  $\mathbf{Z} = p_z \mathbf{P}_{\hat{\mathbf{h}}_b}^\perp$ ,  $\mathbf{P}_{\hat{\mathbf{h}}_b}^\perp = \mathbf{I}_{N_t} - \hat{\mathbf{h}}_b \hat{\mathbf{h}}_b^H / \|\hat{\mathbf{h}}_b\|^2$ , and  $p_z$  is the power invested on AN. We assume that the  $\Phi$  in  $\hat{\mathbf{h}}_b = \mathbf{h}_{ab} + \mathbf{G}_{ar}^H \Phi^H \mathbf{h}_{rb}$  is randomly chosen. The allocated power  $p_w$  and  $p_z$  are optimized in Problem (20) or Problem (47).

- MRT-isoAN-optIRS: The difference of the MRT-isoAN-optIRS scheme from the MRT-isoAN-randIRS scheme is that the utilized  $\Phi$  at the IRS is optimized.
- randIRS: The difference of the Random-IRS scheme from the proposed algorithm is that the utilized  $\Phi$  at the IRS is random.

#### A. The Case of Uncorrelated CSI Errors

Firstly, we consider the uncorrelated CSI errors, where the variance matrix of  $\mathbf{g}_{ge,k}$  is defined as  $\Sigma_{ge,k} = \varepsilon_{g,k}^2 \mathbf{I}$ , and  $\varepsilon_{g,k}^2 = \delta_{g,k}^2 \|\text{vec}(\bar{\mathbf{G}}_{ce,k})\|_2^2$ , while the variance matrix of  $\mathbf{g}_{he,k}$  is defined as  $\Sigma_{he,k} = \varepsilon_{h,k}^2 \mathbf{I}$ , and  $\varepsilon_{h,k}^2 = \delta_{h,k}^2 \|\hat{\mathbf{h}}_{ae,k}\|_2^2$ .  $\delta_{g,k}, \delta_{h,k} \in [0, 1)$  are the normalized CSI errors, which measure the relative amount of CSI errors. Unless specified, the normalized CSI error for partial Eves' CSI error is  $\delta_{g,k} = 0.01$ ,  $\forall k \in \mathcal{K}$ , and the normalized CSI error for full Eves' CSI error is  $\delta_{g,k} = \delta_{h,k} = 0.01$ ,  $\forall k \in \mathcal{K}$  [27].

1) *Convergence Performance*: The convergence performance of the proposed method is investigated in Fig. 3. It is observed that the proposed algorithm monotonically converges with different values of  $N_t$ ,  $M$ , and  $K$  for both the partial and full CSI errors. The convergence speed decreases with a larger number of IRS reflection units  $M$  and a larger number of Eves  $K$ , while increases with a larger number of the transmit antennas  $N_t$ . The number of iterations required for convergence is more sensitive to  $M$  than  $N_t$  and  $K$ . As observed in Fig. 3, the convergence speed with partial CSI errors is generally slower than that with full CSI errors, except the case when  $M$  is relatively large. Since the dimensions of the search space in the optimization subproblems increase with  $N_t$  and  $M$ , and the number of constraints in the optimization subproblems increases with  $K$ , the running time required by the proposed robust design grows with  $N_t$ ,  $M$ , and  $K$ .

2) *Transmit Power vs the Minimum Data Rate of Bob*: Fig. 4 describes the transmit power at different values of  $\log \gamma$ , which is the minimum data rate of Bob. It is observed that the transmit power increases monotonically with  $\log \gamma$ . This means that Alice has to transmit more power to ensure a larger data rate of Bob.

Here, the Eves' data rates can be limited by three factors, which are the CSI errors of Eves' channels, the AN impairment, and the reconfiguration on wireless propagation

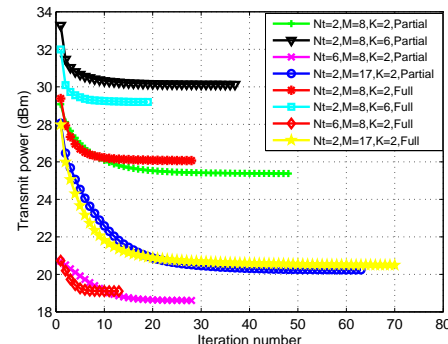


Fig. 3. Convergence performance.

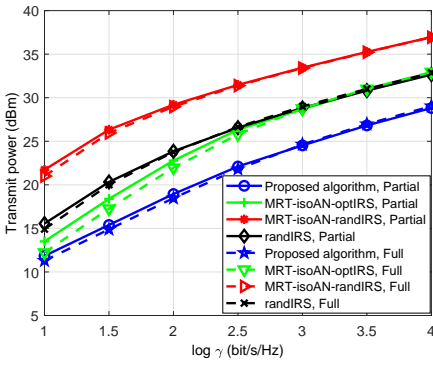


Fig. 4. Required transmit power at each minimum data rate  $\log \gamma$  of Bob, where  $N_t = 2$ ,  $M = 15$ ,  $K = 2$ , and  $\log \beta = 0.5$  bit/s/Hz.

by IRS. For the MRT-isoAN-randIRS and MRT-isoAN-optIRS schemes, the AN is isotropically distributed on the orthogonal complement subspace of the equivalent channel of Bob, thus the AN power directed to Eves is relatively low. For the proposed algorithm and the randIRS scheme, the AN is steered towards Eves' direction, thus the AN power invested on Eves is spatially focused and relatively high.

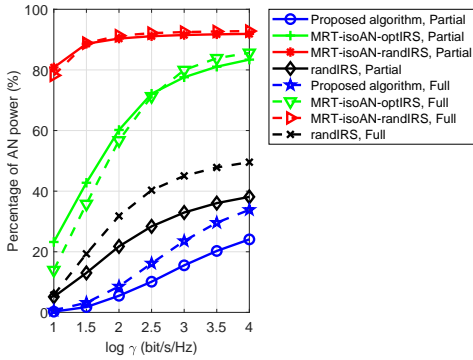


Fig. 5. Percentage of invested AN power at each minimum data rate  $\log \gamma$  of Bob, where  $N_t = 2$ ,  $M = 15$ ,  $K = 2$ , and  $\log \beta = 0.5$  bit/s/Hz.

When  $\log \gamma$  is relatively small, the transmit power required by the MRT-isoAN-optIRS is close to that required by the proposed algorithm, and lower than that of the randIRS scheme. This means that when  $\log \gamma$  is small, the isotropically distributed AN is low but enough to confuse the Eves, and the optimization of the IRS phase shifts seems more effective than the AN focusing. As the  $\log \gamma$  increases, the performance of the MRT-isoAN-optIRS scheme degrades into that of the randIRS scheme, and their required power is higher than that of the proposed algorithm. This signifies that when  $\log \gamma$  is large, the optimization on IRS phase shifts is not enough to interfere Eves, and the AN should be designed and focused on Eves.

It is also observed that when  $\log \gamma$  is small, the transmit power of all these schemes with full CSI errors is smaller than that with partial CSI errors. This is because when  $\log \gamma$  is small, the full CSI errors can impair Eves with larger errors than the partial CSI errors. When  $\log \gamma$  is large, neither the full or the partial CSI errors are enough for impairing Eves, thus more AN power is required. Since it is more difficult to transmit AN through channels with worse quality, the robust design with full CSI errors require more transmit power.

The AN power invested from the total transmit power at different values of  $\log \gamma$  is shown in Fig. 5. As observed, the percentage of AN power increases with  $\log \gamma$  for all schemes. A higher percentage of AN power is required by the MRT-isoAN strategy than the proposed algorithm and randIRS scheme. That's because the AN in the MRT-isoAN strategy is isotropically distributed, while the AN in the proposed algorithm and randIRS scheme is spatially focused.

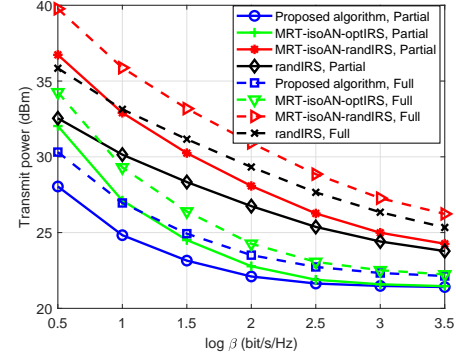


Fig. 6. Required transmit power at each maximum data rate  $\log \beta$  of Eves, where  $N_t = 2$ ,  $M = 15$ ,  $K = 2$ , and  $\log \gamma = 4$  bit/s/Hz.

3) *Transmit Power vs the Maximum Data Rate of Eves:* The transmit power at different values of  $\log \beta$  is depicted in Fig. 6, where  $\log \beta$  denotes the maximum data rate of Eves. As observed, the transmit power monotonically decreases with  $\log \beta$ , and the decreasing speed slows down with increased  $\log \beta$ . Increasing  $\log \beta$  means that the limitation on the information leakage for Eves is relaxed, thus less AN power and transmit power are required. With a relatively large  $\log \beta$ , the transmit power required by the MRT-isoAN-optIRS scheme becomes close to that of the proposed algorithm, and the MRT-isoAN-randIRS scheme becomes close to the randIRS scheme. This means that it is enough to transmit almost isotropically distributed AN by MRT strategy to interfere Eves when the data rate limitation for Eves is relaxed. The transmit power required by the MRT-isoAN-optIRS and the proposed algorithm is much lower than the MRT-isoAN-randIRS and randIRS schemes, which demonstrates the necessity to optimize the phase shifts of the IRS. It is also found that the transmit power with full CSI errors is higher than that with partial CSI errors, and the gap is narrowed gradually with  $\log \beta$ .

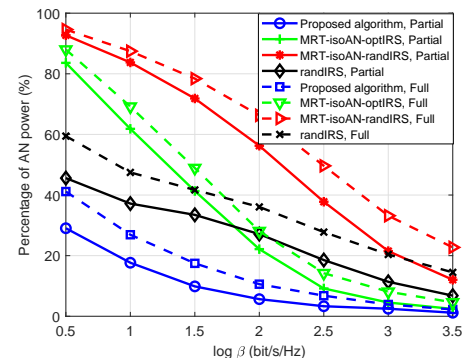


Fig. 7. Percentage of invested AN power at each maximum data rate  $\log \beta$  of Eves, where  $N_t = 2$ ,  $M = 15$ ,  $K = 2$ , and  $\log \gamma = 4$  bit/s/Hz.



The invested AN power from the total transmit power at different values of  $\log \beta$  is shown in Fig. 7. As observed, the AN power decreases with  $\log \beta$ . The tighter the limitation for information leakage is, the more AN power is required. The AN power of the proposed algorithm and Random-IRS scheme is less than that of the Optimized-MRT and Random-MRT, which means that the idea of making AN spatially focused is beneficial for reducing the transmit power.

#### 4) The Impact of the Number of IRS Reflection Units:

Fig. 8 investigates the impact of the unit number of IRS on reducing the transmit power. It is observed that the transmit power can be reduced with more IRS units for all these schemes. The robust transmission design of the proposed algorithm outperforms the other schemes in terms of the lowest transmit power. The transmit power consumption of the proposed algorithm and the MRT-isoAN-optIRS scheme drops much faster than that of the randIRS and MRT-isoAN-randIRS schemes, since the phase shifts of the IRS are optimized for them. When  $M$  becomes large, the power consumption of the MRT-isoAN-optIRS tends to approach that of the proposed algorithm. This is because larger  $M$  brings more cascaded CSI errors, which enhance the capability of impairing Eves by the isoAN strategy. The transmit power with the full CSI errors is slightly larger than that with the partial CSI errors.

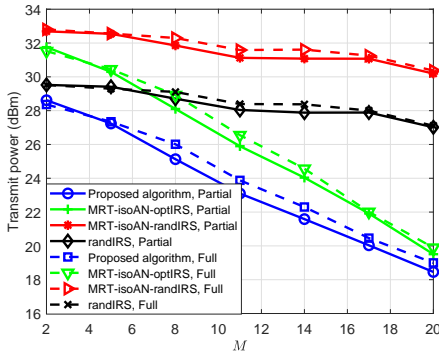


Fig. 8. Required transmit power with different unit numbers  $M$  of the IRS, where  $N_t = 2$ ,  $K = 2$ ,  $\log \beta = 1$  bit/s/Hz, and  $\log \gamma = 3$  bit/s/Hz.

5) *The Impact of the Number of Transmit Antennas:* Fig. 9 investigates the impact of the number of Alice's antennas  $N_t$  on the transmit power consumption. As observed, the transmit power decreases with  $N_t$  for all schemes, which means that more transmit antennas can reduce the transmit power in the robust design. The proposed algorithm and the randIRS scheme are more sensitive to  $N_t$  than the MRT-isoAN-optIRS scheme and MRT-isoAN-randIRS scheme. This is because a larger  $N_t$  is helpful for increasing the Bob's data rate, while is not conducive to decrease the Eves' data rate. The AN power for the isoAN strategy is almost isotropically distributed, which is less effective to reduce the transmit power when  $N_t$  is large. Thus, the transmit power required by the MRT-isoAN-optIRS and MRT-isoAN-randIRS schemes reduces gently. The transmit power with full CSI errors is slightly higher than that with partial CSI errors.

6) *The Impact of Channel Errors:* The impact of channel uncertainty on the transmit power is investigated in Fig. 10, where the channel uncertainty is measured by the normalized

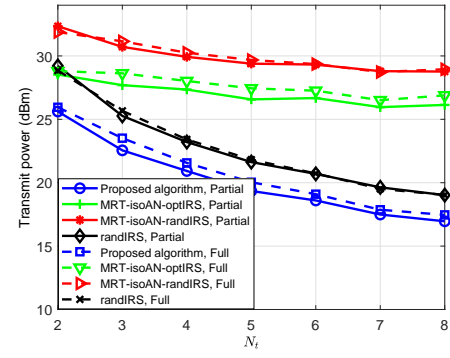


Fig. 9. Required transmit power with different transmit antenna numbers  $N_t$ , where  $M = 8$ ,  $K = 2$ ,  $\log \beta = 1$  bit/s/Hz, and  $\log \gamma = 3$  bit/s/Hz.

CSI error  $\delta$  of Eves' channels, and we assume that  $\delta_{g,k} = \delta_{h,k} = \delta$ ,  $\forall k$ . It is found that the transmit power consumption increases with the normalized CSI errors  $\delta$ , which means that more power has to be transmitted when the channel quality degrades. Under a relatively large CSI error, the transmit power with full CSI errors is slightly lower than that with partial CSI errors. This means that when the CSI error is large, the full CSI errors are preferred to impair the Eves. It is also observed that when full CSI errors exist for Eves, the transmit power changing with the channel error of the direct link is higher than that changing with the channel error of the reflecting link. This signifies that the robust transmission design is more sensitive to the CSI error of the direct link. Among all these schemes, the proposed algorithm requires the lowest transmit power. By comparing with the No-IRS scheme, it demonstrates that the security performance of the IRS-aided communication system is superior to the no-IRS system even with CSI errors of both the direct link and IRS reflecting link.

7) *The Impact of the Number of Eves:* Fig. 11 investigates the impact of the number of Eves on transmit power. As observed, the transmit power increases with the number of Eves for all schemes. When the number of Eves is relatively small, the transmit power with partial CSI errors is slightly lower than that with full CSI errors, which means that the partial CSI error is sufficient to impair Eves. When more Eves exist, the transmit power with partial CSI errors is slightly higher than that with full CSI errors, which means the full CSI error is preferred to interfere more Eves.

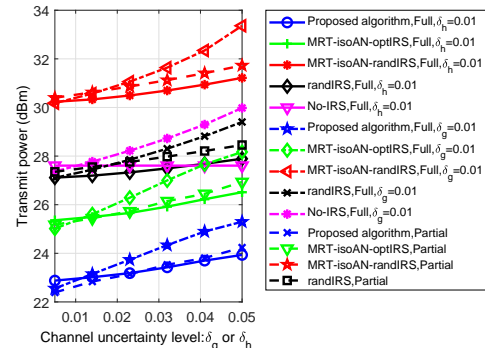


Fig. 10. Required transmit power with different channel uncertainty levels  $\delta$ , where  $N_t = 2$ ,  $M = 12$ ,  $K = 3$ ,  $\log \beta = 1$  bit/s/Hz, and  $\log \gamma = 3$  bit/s/Hz.

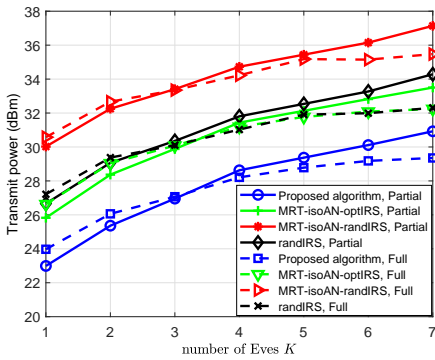


Fig. 11. Required transmit power with different numbers of Eves  $K$ , where  $N_t = 2$ ,  $M = 8$ ,  $\log \beta = 1$  bit/s/Hz, and  $\log \gamma = 3$  bit/s/Hz.

8) *The Impact of the Initial Values on the Performance of the Proposed Algorithm:* Due to the nonconvexity of Problem (10), different initial values may result in different locally optimal solutions obtained by the proposed algorithm. It is difficult to find a good initial value for  $\phi$  since a good  $\phi$  should enhance the Alice-Bob channel and deteriorate the Alice-Eve channel. Thus, we choose initial  $\phi$  randomly, but generate the initial  $\mathbf{w}$  and  $\mathbf{Z}$  by solving Problem (20) or Problem (47). To study the impact of the initialization of  $\mathbf{w}$  and  $\mathbf{Z}$  on the performance of the proposed algorithm, we test 30 randomly generated channels shown in Fig. 12, where AO-OPT refers to using optimized  $\mathbf{w}$  and  $\mathbf{Z}$  as initial values and AO-EXH refers to that  $\mathbf{w}$  and  $\mathbf{Z}$  are initialized by the best of 1000 feasible random normalized vectors and matrices for each channel realization. It can be seen that the total power consumption of AO-OPT is almost the same as that of AO-EXH, implying that the optimized  $\mathbf{w}$  and  $\mathbf{Z}$  is a good option for the initialization.

9) *The Multiple-Bob Case:* We further investigate the effectiveness of the proposed robust transmission design in the multiple-Bob case. It is noted that the baseline scheme of MRT and isotropic AN was proposed aiming for the single-Bob case [35]. If we extend this baseline scheme into the multiple-Bob case directly, the  $k$ th beamformer becomes  $\mathbf{w}_k = \sqrt{p_k} \frac{\hat{\mathbf{h}}_{b,k}}{\|\hat{\mathbf{h}}_{b,k}\|}$ , and the AN covariance matrix becomes  $\mathbf{Z} = p_z \mathbf{P}_{\hat{\mathbf{H}}_b}^\perp$ ,  $\mathbf{P}_{\hat{\mathbf{H}}_b}^\perp = \mathbf{I}_{N_t} - \check{\mathbf{H}}_b \check{\mathbf{H}}_b^H / \|\check{\mathbf{H}}_b\|_F^2$ ,  $\check{\mathbf{H}}_b = [\hat{\mathbf{h}}_{b,1}, \hat{\mathbf{h}}_{b,2}, \dots, \hat{\mathbf{h}}_{b,K}]$ , where the  $p_k$  and  $p_z$  are power scale factors to be optimized. However, during the simulations, we find that the feasibility rate of this kind of baseline schemes decreases greatly, and becomes zero

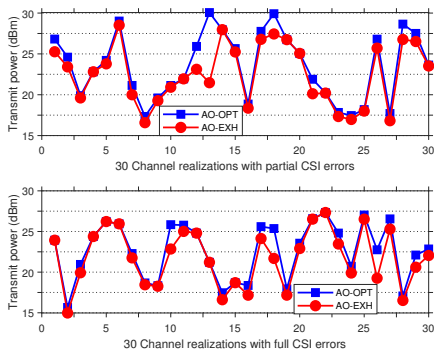


Fig. 12. Performance comparison of AO-OPT and AO-EXH, where  $N_t = 2$ ,  $M = 15$ ,  $K = 2$ ,  $\log \gamma = 2.5$  bit/s/Hz, and  $\log \beta = 0.5$  bit/s/Hz.

when the required secrecy rate is slightly high. The reason is that this form of  $\mathbf{w}_k$  and  $\mathbf{Z}$  makes the data rate of each Bob stay at a low level and cannot ensure a good secrecy rate. This means that the baseline schemes of MRT and isotropic AN cannot guarantee the system security when multiple Bobs coexist with multiple Eves.

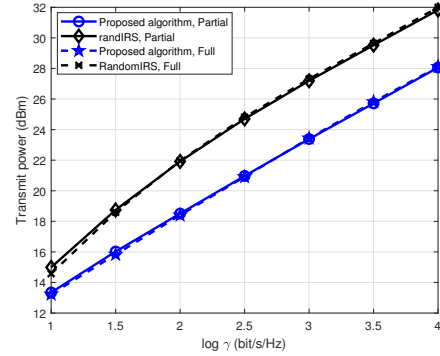


Fig. 13. Required transmit power at each minimum data rate  $\log \gamma$  of multiple Bobs, where  $N_t = 2$ ,  $M = 15$ ,  $K = 2$ ,  $\log \beta = 0.5$  bit/s/Hz,  $\gamma_k = \gamma$ ,  $\beta_k = \beta$ ,  $\forall k \in \mathcal{K}$ .

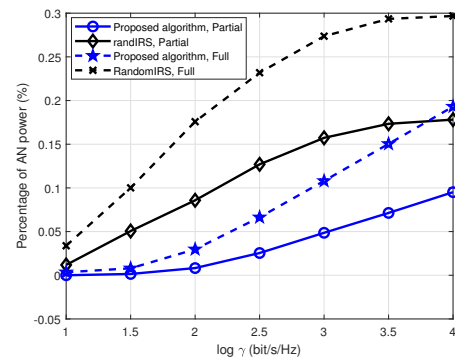


Fig. 14. Percentage of invested AN power at each minimum data rate  $\log \gamma$  of multiple Bobs, where  $N_t = 2$ ,  $M = 15$ ,  $K = 2$ ,  $\log \beta = 0.5$  bit/s/Hz,  $\gamma_k = \gamma$ ,  $\beta_k = \beta$ ,  $\forall k \in \mathcal{K}$ .

Thus, to make the comparison fair, we only compare our proposed algorithm with the randIRS scheme. Fig. 13 describes the transmit power of different schemes at different values of  $\log \gamma$  in the multiple-Bob case, where two Bobs are located at (3,36,0) m and (3,48,0) m. The percentage of AN power invested from the total transmit power at different values of  $\log \gamma$  is shown in Fig. 14. It is observed from Fig. 13 that the transmit power of the proposed algorithm is much lower than the randIRS scheme. From Fig. 14, we find that the percentage of AN power increases with  $\log \gamma$ , and invested AN power of the proposed algorithm is much slower than the randIRS scheme. Moreover, as compared with the single-Bob case, the required AN power is greatly reduced in the multi-Bob case and is almost negligible.

### B. The Case of Correlated CSI Errors

In this subsection, we consider the robust transmission design with correlated CSI errors, where  $\Sigma_{ge,k}^{1/2} = \varepsilon_{g,k} \mathbf{S}_{MN_t}$ ,  $\Sigma_{he,k}^{1/2} = \varepsilon_{h,k} \mathbf{S}_{N_t}$ ,  $[\mathbf{S}_{MN_t}]_{m,n}$ ,  $[\mathbf{S}_{N_t}]_{m,n} = \eta^{|m-n|}$ , and  $\eta$  is set as 0.9 [40]. To make the relative amount of correlated CSI errors comparable with the case of uncorrelated CSI errors, we

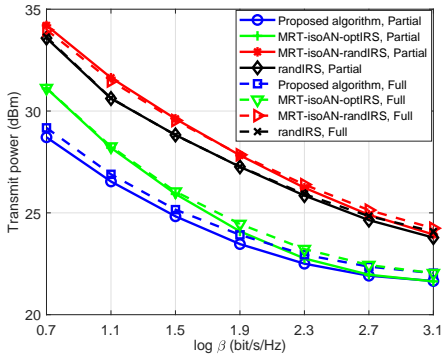


Fig. 15. Required transmit power at each  $\log\beta$  with correlated CSI errors, where  $N_t = 2$ ,  $M = 15$ ,  $K = 2$ , and  $\log\gamma = 3.5$  bit/s/Hz.

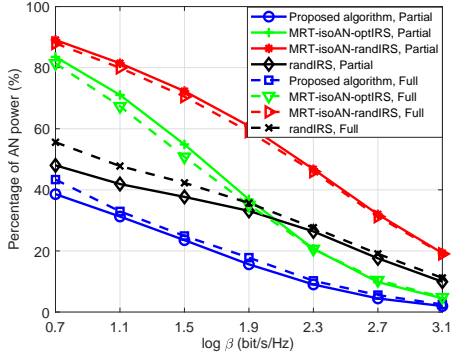


Fig. 16. Percentage of invested AN power at each  $\log\beta$  with correlated CSI errors, where  $N_t = 2$ ,  $M = 15$ ,  $K = 2$ , and  $\log\gamma = 3.5$  bit/s/Hz.

assume that  $\varepsilon_{g,k}^2 = \delta_{g,k}^2 \|\text{vec}(\bar{\mathbf{G}}_{ce,k})\|_2^2 [\|\mathbf{I}_{MN_t}\|_F / \|\mathbf{S}_{MN_t}^2\|_F]$ , and  $\varepsilon_{h,k}^2 = \delta_{h,k}^2 \|\bar{\mathbf{h}}_{ae,k}\|_2^2 [\|\mathbf{I}_{N_t}\|_F / \|\mathbf{S}_{N_t}^2\|_F]$ .

The transmit power versus the maximum tolerable channel capacity  $\log\beta$  of Eves is shown in Fig. 15. It is shown that the transmit power of the proposed algorithm is the lowest among all the schemes, which demonstrates the effectiveness and superiority of the proposed algorithm for the uncorrelated CSI errors. The gap between the proposed algorithm and the Optimized-MRT scheme becomes smaller as  $\log\beta$  increases. The transmit power of the full CSI errors is slightly higher than that of the partial CSI errors.

The invested AN power from the total transmit power with different  $\log\beta$  is shown in Fig. 16. The required AN power reduces when the data rate limit of Eves is relaxed. The AN power with the full CSI errors is slightly higher than that with the partial CSI errors for the proposed algorithm.

## V. CONCLUSION

An robust transmission strategy was designed for an IRS-aided secure communication system by considering the statistical CSI errors on both direct links and IRS reflecting links to Eves. We proposed an AO algorithm to solve the formulated OCR-PM problem by leveraging the BTI, SDR technique, and CCP method. The proposed algorithm can apply with the uncorrelated and correlated CSI errors, and its superiority over baseline schemes is verified by simulations.

## APPENDIX A PROOF OF LEMMA 2

We define a matrix  $\mathbf{C} = \mathbf{a}\mathbf{b}^T$ , where  $\mathbf{a}$  and  $\mathbf{b}$  are vectors. Since we have  $\text{vec}(\mathbf{C}) = \mathbf{b} \otimes \mathbf{a}$ , then we have

$$\begin{aligned} \|\mathbf{C}\|_F &= \|\mathbf{b} \otimes \mathbf{a}\|_2 = \sqrt{\text{Tr}(\mathbf{C}\mathbf{C}^H)} = \sqrt{\text{Tr}(\mathbf{a}\mathbf{b}^T\mathbf{b}^*\mathbf{a}^H)} \\ &= \sqrt{\mathbf{b}^T\mathbf{b}^*\mathbf{a}^H\mathbf{a}} = \|\mathbf{b}\|_2 \|\mathbf{a}\|_2. \end{aligned} \quad (73)$$

## APPENDIX B PROOF OF (32d)

By utilizing the rank properties  $\text{rank}\{\mathbf{A}\mathbf{B}\} \leq \min\{\text{rank}(\mathbf{A}), \text{rank}(\mathbf{B})\}$  and  $\text{rank}(\mathbf{A} \otimes \mathbf{B}) = \text{rank}(\mathbf{A}) \cdot \text{rank}(\mathbf{B})$ , we can prove  $\lambda_{\min}\{\boldsymbol{\Sigma}_{e,k}^{1/2}((c_g\mathbf{I}_{N_t})^T \otimes \mathbf{E})\boldsymbol{\Sigma}_{e,k}^{1/2}\} = 0$  as follows.

We have the inequality of  $\text{rank}\{\boldsymbol{\Sigma}_{ge,k}^{1/2}((c_g\mathbf{I}_{N_t})^T \otimes \mathbf{E})\boldsymbol{\Sigma}_{ge,k}^{1/2}\} \leq \min\{\text{rank}(\boldsymbol{\Sigma}_{ge,k}^{1/2}), \text{rank}((c_g\mathbf{I}_{N_t})^T \otimes \mathbf{E})\}$ . Since  $\text{rank}((c_g\mathbf{I}_{N_t})^T \otimes \mathbf{E}) = \text{rank}(c_g\mathbf{I}_{N_t}) \cdot \text{rank}(\mathbf{E}) = N_t$ , we have  $\text{rank}\{\boldsymbol{\Sigma}_{ge,k}^{1/2}((c_g\mathbf{I}_{N_t})^T \otimes \mathbf{E})\boldsymbol{\Sigma}_{ge,k}^{1/2}\} \leq N_t$ . Since the matrix  $\boldsymbol{\Sigma}_{ge,k}^{1/2}((c_g\mathbf{I}_{N_t})^T \otimes \mathbf{E})\boldsymbol{\Sigma}_{ge,k}^{1/2}$  is an  $MN_t \times MN_t$  dimensional semidefinite matrix, its minimum eigenvalue is zero as

$$\lambda_{\min}\{\boldsymbol{\Sigma}_{ge,k}^{1/2}((c_g\mathbf{I}_{N_t})^T \otimes \mathbf{E})\boldsymbol{\Sigma}_{ge,k}^{1/2}\} = 0. \quad (74)$$

## APPENDIX C PROOF OF LEMMA 3

For any three matrices  $\mathbf{O}$ ,  $\mathbf{P}$ ,  $\mathbf{Q}$ , we have

$$\begin{aligned} \|\mathbf{P}\mathbf{O}\mathbf{Q}\|^2 &\stackrel{(w_1)}{=} \text{Tr}[\mathbf{P}\mathbf{O}\mathbf{Q}\mathbf{Q}^H\mathbf{O}^H\mathbf{P}^H] = \text{Tr}[\mathbf{O}^H\mathbf{P}^H\mathbf{P}\mathbf{O}\mathbf{Q}\mathbf{Q}^H] \\ &\stackrel{(w_2)}{=} \text{vec}^H(\mathbf{O})[(\mathbf{Q}\mathbf{Q}^H)^T \otimes (\mathbf{P}^H\mathbf{P})]\text{vec}(\mathbf{O}) \\ &= \text{vec}^H(\mathbf{O})[(\mathbf{Q}^T \otimes \mathbf{P})^H(\mathbf{Q}^T \otimes \mathbf{P})]\text{vec}(\mathbf{O}) \\ &= \|(\mathbf{Q}^T \otimes \mathbf{P})\text{vec}(\mathbf{O})\|^2, \end{aligned} \quad (75)$$

where  $(w_1)$  is due to  $\|\mathbf{A}\|_F^2 = \text{Tr}[\mathbf{A}\mathbf{A}^H]$ , and  $(w_2)$  is obtained by invoking the identity  $\text{Tr}(\mathbf{A}^H\mathbf{B}\mathbf{C}\mathbf{D}) = \text{vec}^H(\mathbf{A})(\mathbf{D}^T \otimes \mathbf{B})\text{vec}(\mathbf{C})$ .

## APPENDIX D PROOF OF THAT $\mathbf{B}_{e,k}$ IS RANK DEFICIENT

According to the definition of  $\mathbf{B}_{e,k}$ , it can be equivalently expressed as  $\mathbf{B}_{e,k} = \mathbf{C}_{e,k}^H \mathbf{C}_{e,k}$ , where  $\mathbf{C}_{e,k}$  is defined by

$$\mathbf{C}_{e,k} = \begin{bmatrix} c_g^{1/2} \mathbf{I}_{N_t} \boldsymbol{\Sigma}_{he,k}^{1/2} & (c_g^{1/2} \mathbf{I}_{N_t} \otimes \phi^H) \boldsymbol{\Sigma}_{ge,k}^{1/2*} \end{bmatrix} \quad (76)$$

Then we have  $\lambda_{\min}\{\mathbf{B}_{e,k}\} = \lambda_{\min}\{\mathbf{C}_{e,k}^H \mathbf{C}_{e,k}\}$ . We can find

$$\begin{aligned} \text{rank}\{\mathbf{B}_{e,k}\} &\leq \text{rank}\{\mathbf{C}_{e,k}\} \\ &\leq \text{rank}\left\{ \begin{bmatrix} c_g^{1/2} \mathbf{I}_{N_t} \boldsymbol{\Sigma}_{he,k}^{1/2} & (c_g^{1/2} \mathbf{I}_{N_t} \otimes \phi^H) \boldsymbol{\Sigma}_{ge,k}^{1/2*} \end{bmatrix} \right\} \\ &\leq \text{rank}\{c_g^{1/2} \mathbf{I}_{N_t} \boldsymbol{\Sigma}_{he,k}^{1/2}\} + \text{rank}\{(c_g^{1/2} \mathbf{I}_{N_t} \otimes \phi^H) \boldsymbol{\Sigma}_{ge,k}^{1/2*}\} \leq 2N_t, \end{aligned} \quad (77)$$

where the properties  $\text{rank}(\mathbf{A}\mathbf{B}) \leq \min\{\text{rank}(\mathbf{A}), \text{rank}(\mathbf{B})\}$  and  $\text{rank}([\mathbf{A}, \mathbf{B}]) \leq \text{rank}(\mathbf{A}) + \text{rank}(\mathbf{B})$  are utilized. Since the matrix  $\mathbf{B}_{e,k}$  is a  $(MN_t + N_t) \times (MN_t + N_t)$  dimensional semidefinite matrix, its minimum eigenvalue is zero as  $\lambda_{\min}\{\mathbf{B}_{e,k}\} = 0$  for any  $M \geq 2$ .



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